

Designing A Low-Latency Cuckoo Hash Table for Write-Intensive Workloads Using RDMA

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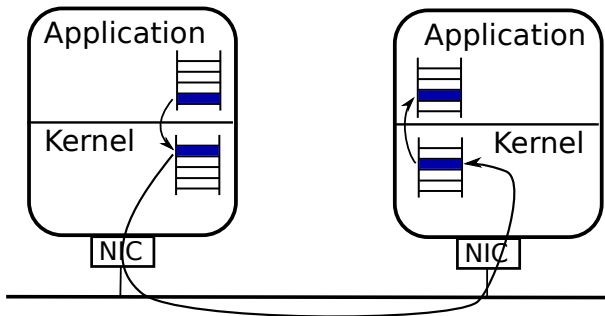
Cheriton School of Computer Science
University of Waterloo

April 13, 2014

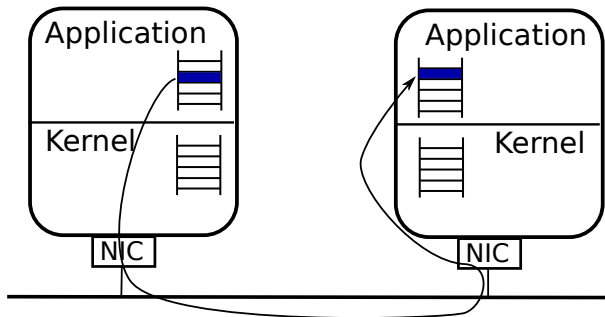
In-Memory Distributed Systems



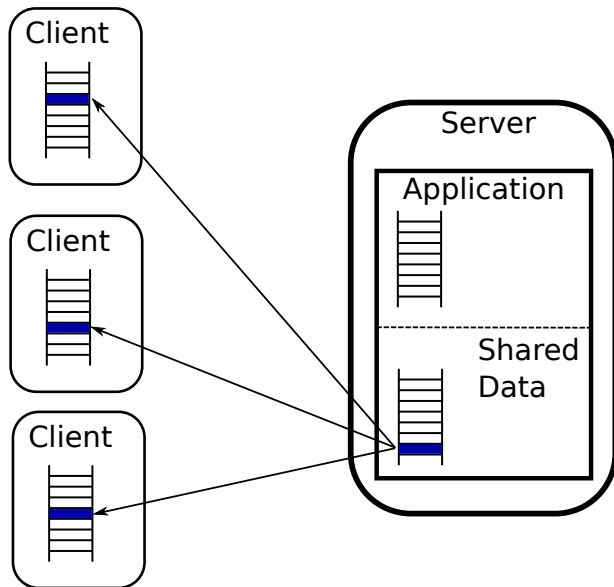
TCP Data Transfer



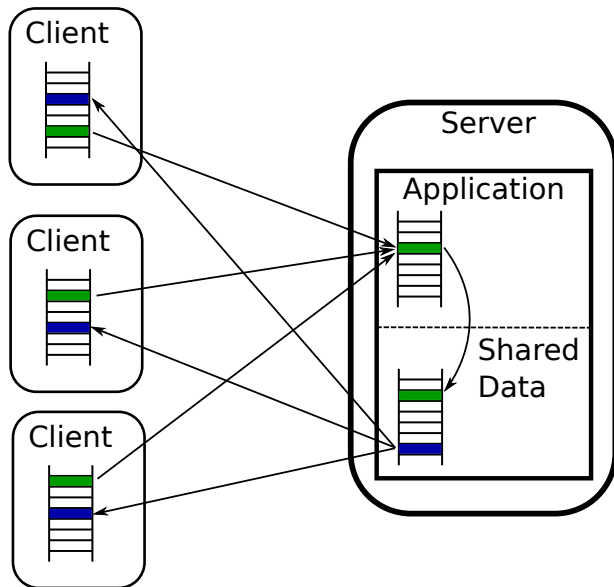
RDMA Data Transfer



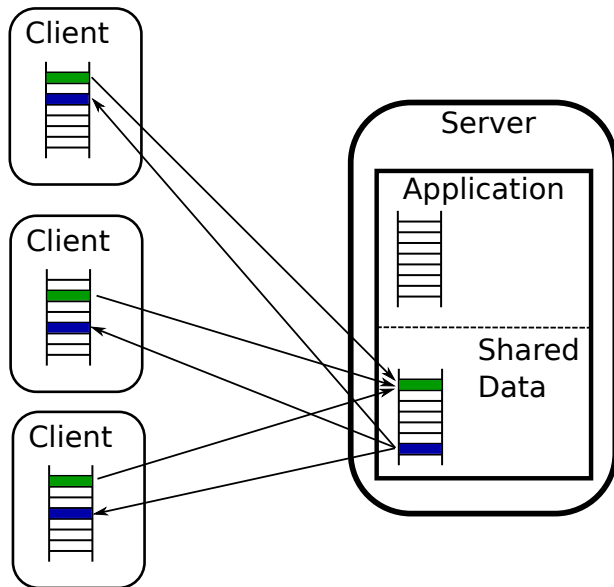
Current State of the Art: Pilaf/FaRM



Current State of the Art: Pilaf/FaRM



One-Sided RDMA for Reads and Writes



- ▶ Cuckoo hash table where clients manage both reads and writes
- ▶ Design:
 - ▶ One-sided RDMA for reads and writes of key-value pairs
 - ▶ Checksum to detect corrupted reads
 - ▶ RDMA compare-and-swap to create atomic writes
 - ▶ Lock-free to eliminate the possibility of deadlock
- ▶ Advantage: Full benefits from one-sided RDMA
- ▶ Tradeoffs compared to the state of the art:
 - ▶ More roundtrips per operation
 - ▶ Added complexity

Cuckoo Hash Table

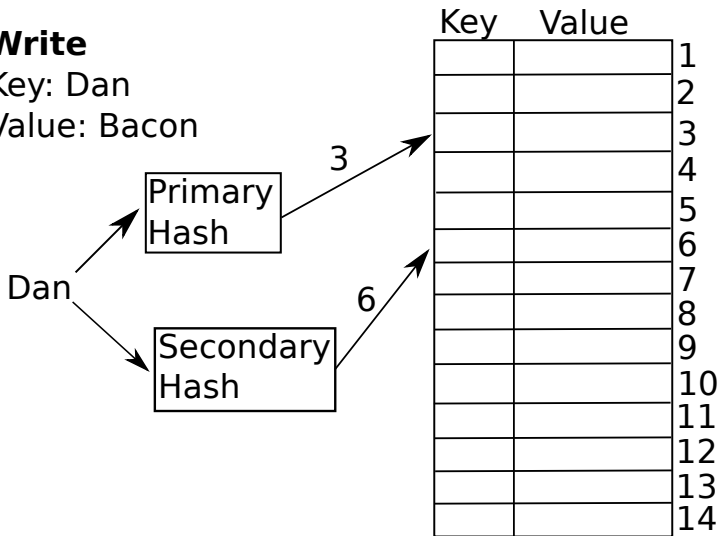
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		1
		2
		3
		4
		5
		6
		7
		8
		9
		10
		11
		12
		13
		14

Cuckoo Hash Table

Write

Key: Dan

Value: Bacon

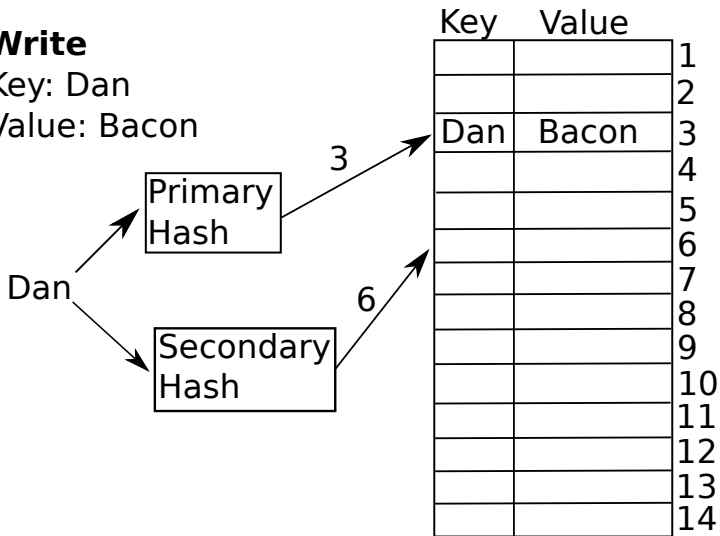


Cuckoo Hash Table

Write

Key: Dan

Value: Bacon

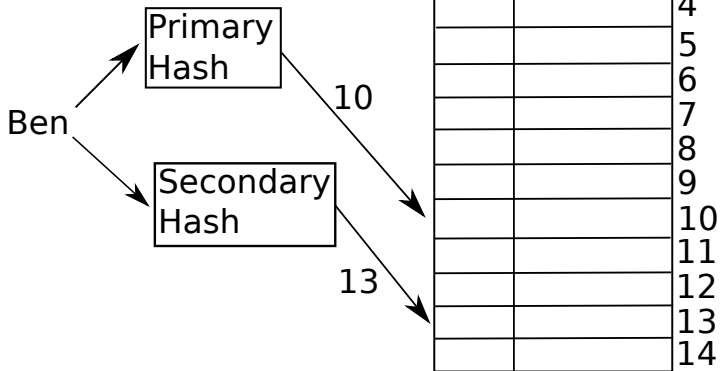


Cuckoo Hash Table

Write

Key: Ben

Value: Apple

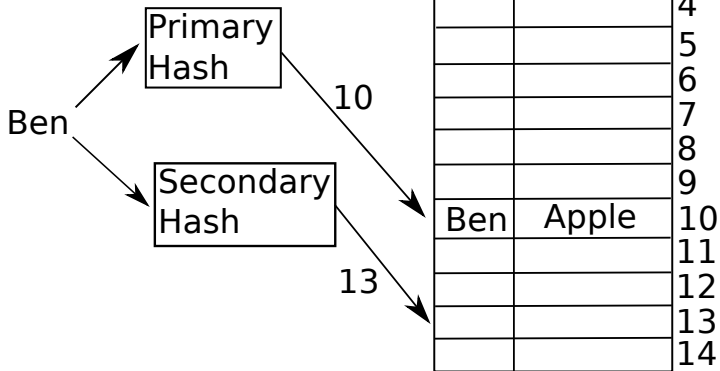


Cuckoo Hash Table

Write

Key: Ben

Value: Apple

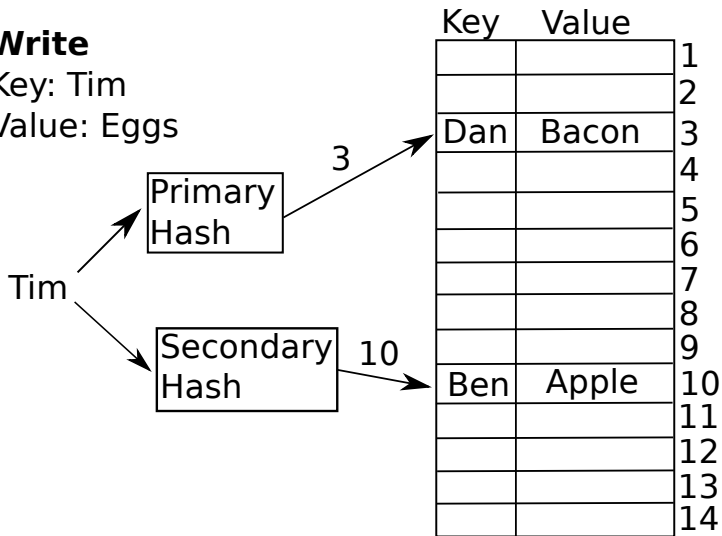


Cuckoo Hash Table

Write

Key: Tim

Value: Eggs

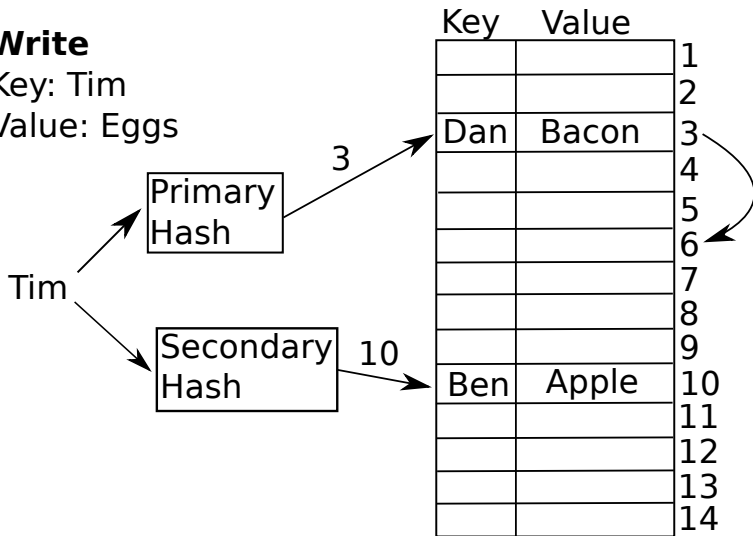


Cuckoo Hash Table

Write

Key: Tim

Value: Eggs

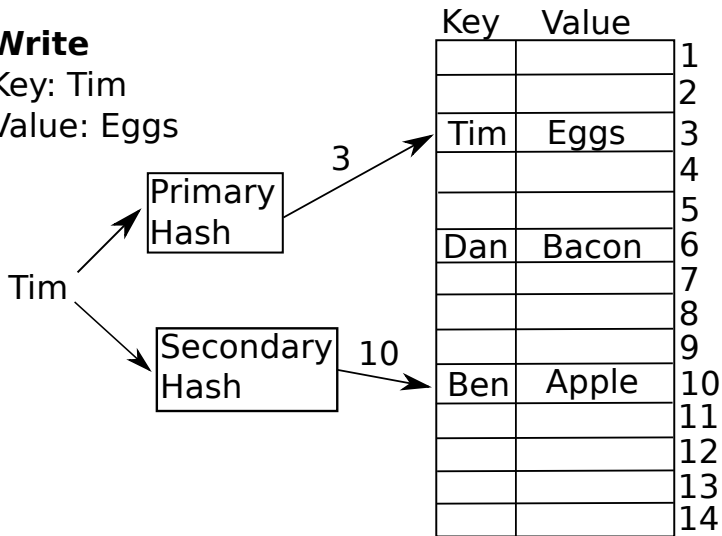


Cuckoo Hash Table

Write

Key: Tim

Value: Eggs



Index Table

Data Table

Index Table

Index	V#

← 64 Bits →

Data Table

Index Table

Index	V#

← 64 Bits →

Data Table

Index Table

Index	V#

← 64 Bits →

Data Table
Key Value Valid C

				1
				2
				3
				4
				5
				6
				7
				8
				9
				10
				11
				12
				13
				14

← N Bytes →

Index Table

Index	V#

← 64 Bits →

Data Table
Key Value Valid C

				1	
				2	
				3	Client 1
				4	
				5	
				6	
				7	Client 2
				8	
				9	
				10	
				11	
				12	Client 3
				13	
				14	

← N Bytes →

**Index
Table**

Data Table

Key	Value	Valid	
			1
			2
			3
			4
			5
			6
			7
			8
			9
			10
			11
			12
			13
			14

Nessie: Write

Write

Key: Tim

Value: Eggs

Index Table

10
6
1
5

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
			7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write

Write Data

Write

Key: Tim

Value: Eggs

Index Table

10
6
1
5

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Read Index Table

Write

Key: Tim

Value: Eggs

Index Table

	10
	6
	1
	5

Tim

Primary →

Secondary →

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

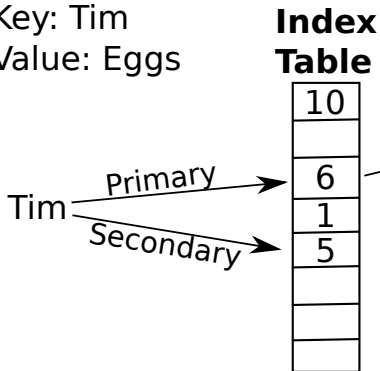
Nessie: Write

Read Data at Primary

Write

Key: Tim

Value: Eggs



Data Table

Key Value Valid

Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

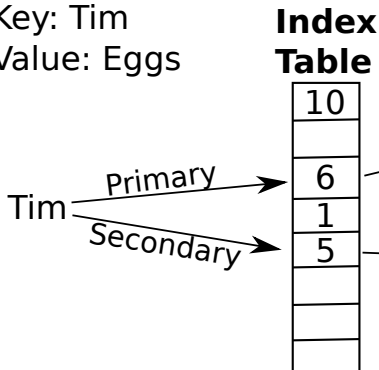
Nessie: Write

Read Data at Secondary

Write

Key: Tim

Value: Eggs



Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

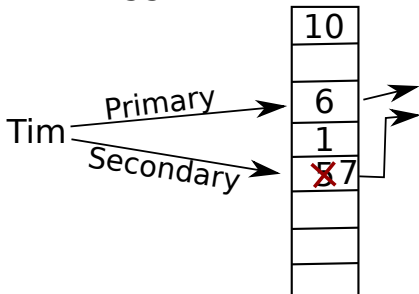
Update Index Table

Write

Key: Tim

Value: Eggs

Index Table



Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

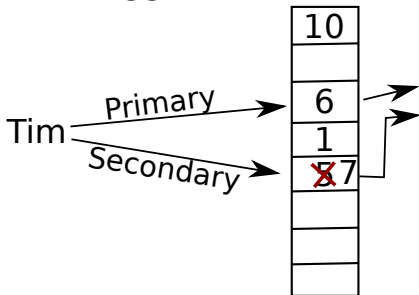
Make Data Valid

Write

Key: Tim

Value: Eggs

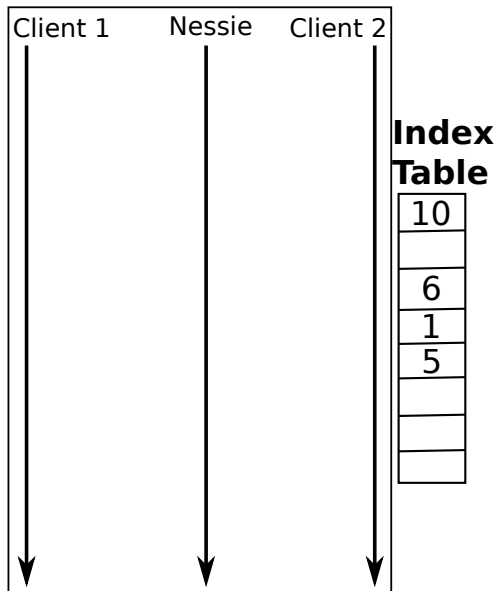
Index Table



Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	T	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

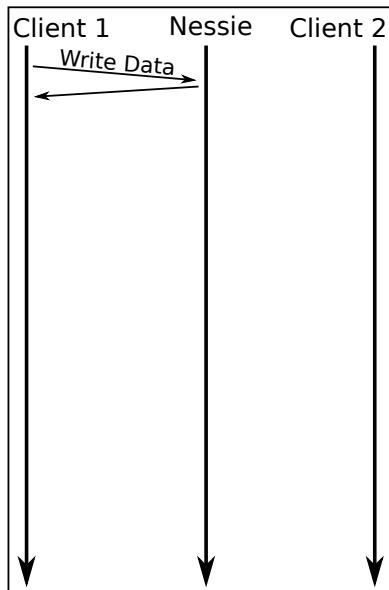
Nessie: Write with Conflict



Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
			7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



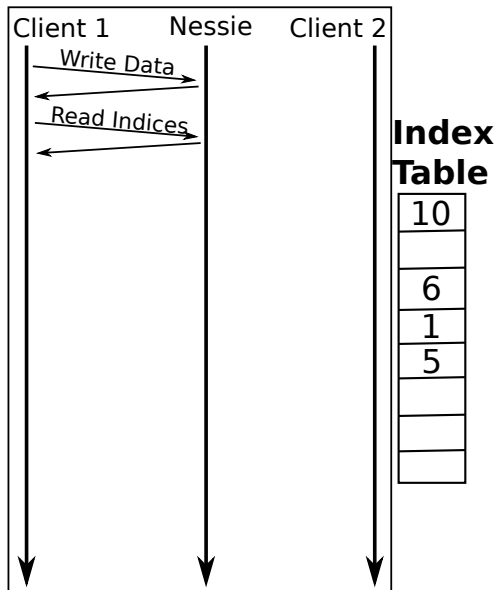
Index Table

10
6
1
5

Data Table

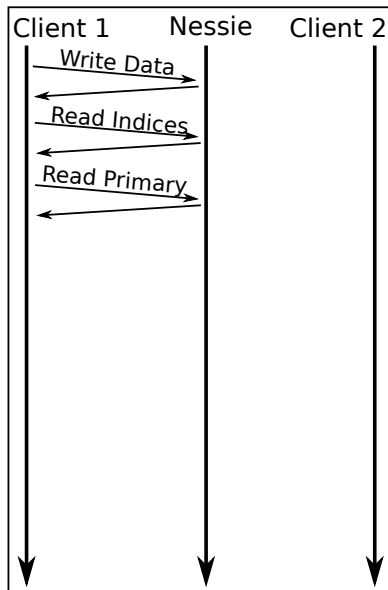
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Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



Data Table			
Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



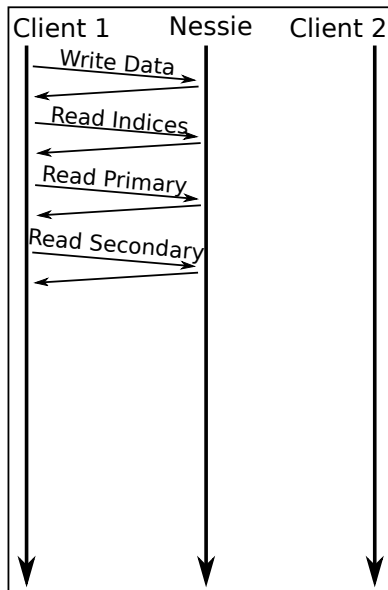
Index Table

10
6
1
5

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



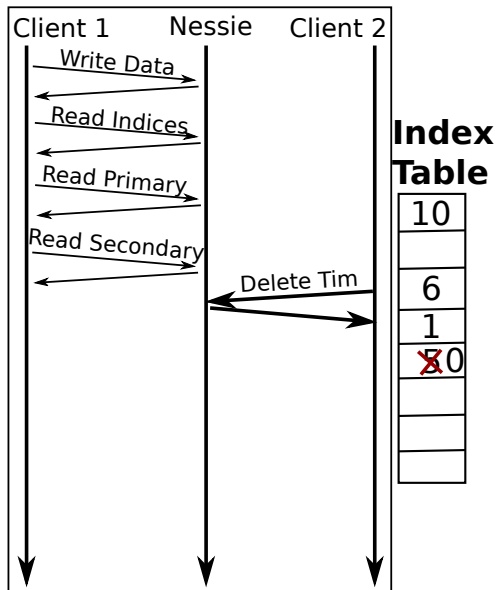
Index Table

10
6
1
5

Data Table

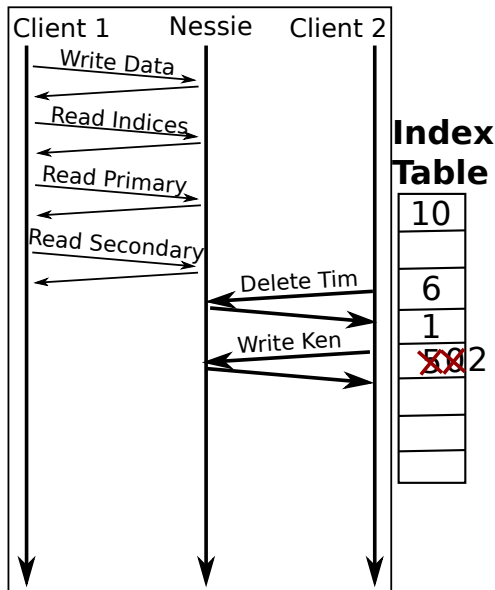
Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



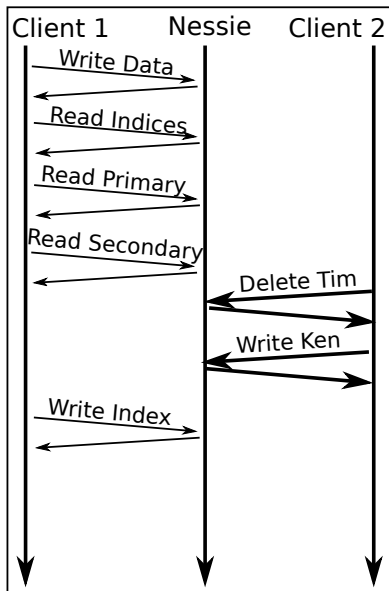
Data Table			
Key	Value	Valid	
Sam	Pizza	T	1
			2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



Data Table			
Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



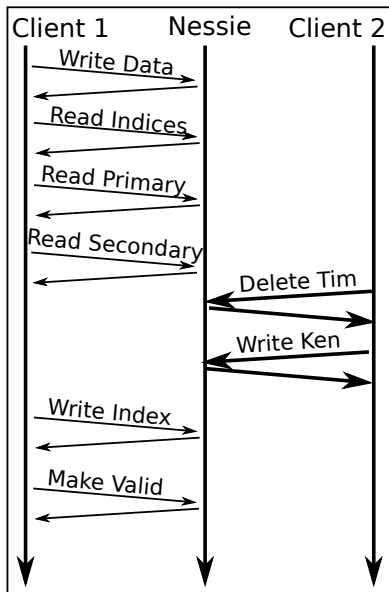
Index Table

10
6
1
50127

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Write with Conflict



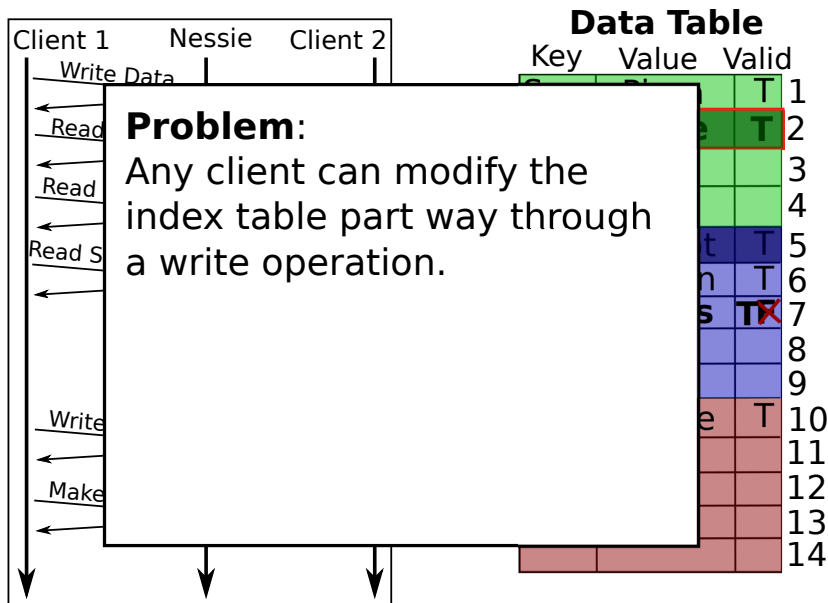
Index Table

10
6
1
5
2
7

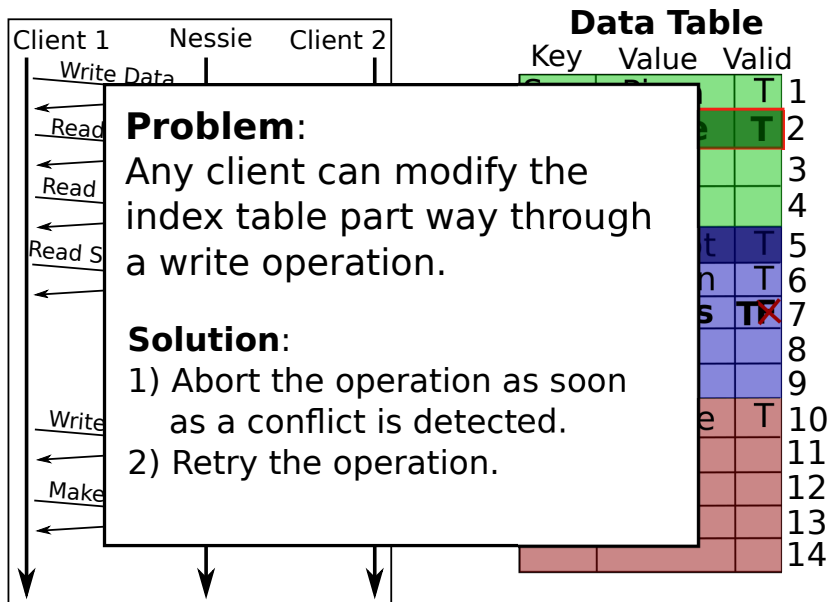
Data Table

Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	T	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

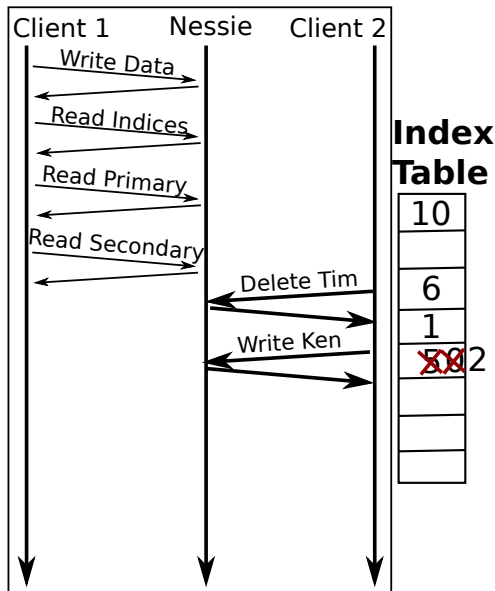
Nessie: Write with Conflict



Nessie: Write with Conflict

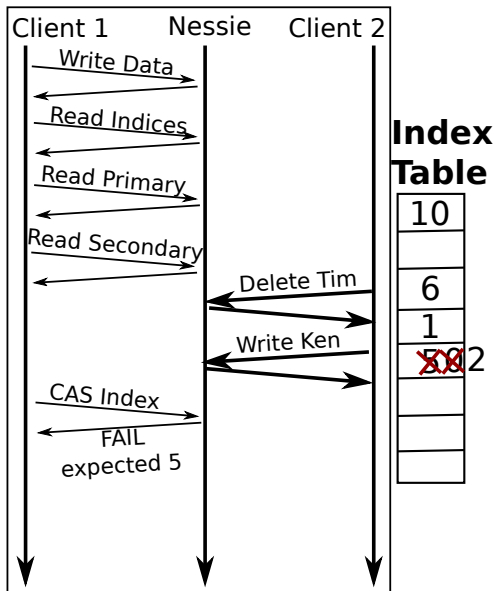


Nessie: Handling Conflicts



Data Table			
Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Handling Conflicts



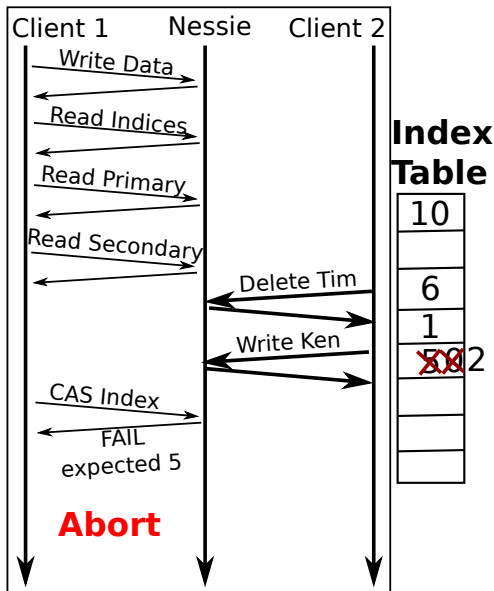
Index Table

10
6
1
5 2

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Handling Conflicts



Index Table

10
6
1
5 2

Data Table

Key	Value	Valid	
Sam	Pizza	T	1
Ken	Cake	T	2
			3
			4
Tim	Carrot	T	5
Dan	Bacon	T	6
Tim	Eggs	F	7
			8
			9
Ben	Apple	T	10
			11
			12
			13
			14

Nessie: Migration

Migrate

Key: Dan

Destination →

Source →

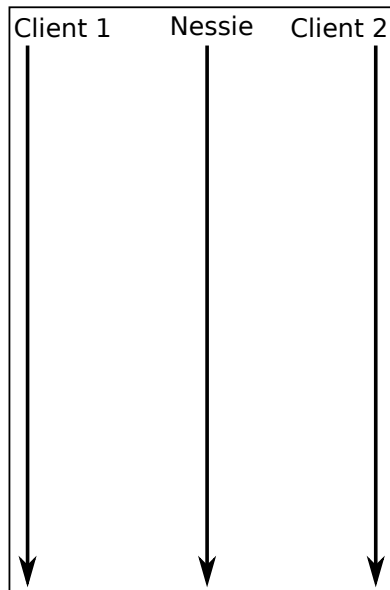
Index Table

10
1
5
11

Data Table

Key	Value	Valid	
Dan	Bacon	T	1
			2
			3
			4
Sam	Pizza	T	5
			6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



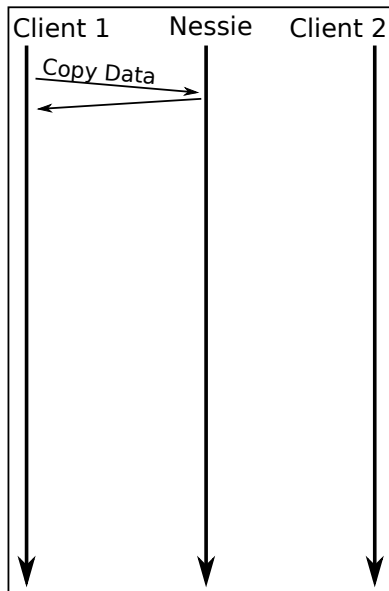
Index Table

10
1
5
11

Data Table

Key	Value	Valid	
Dan	Bacon	T	1
			2
			3
			4
Sam	Pizza	T	5
			6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



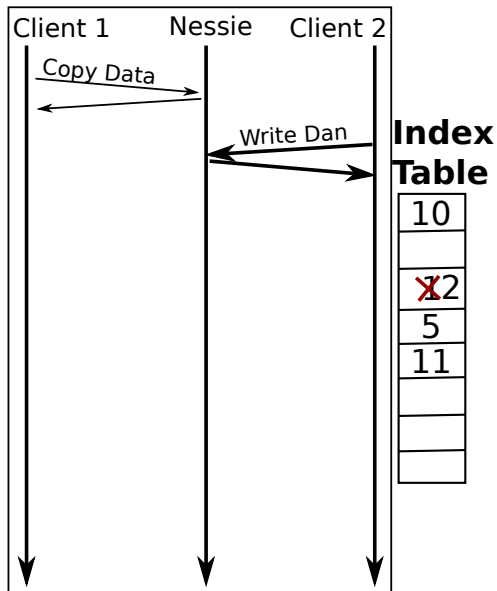
Index Table

10
1
5
11

Data Table

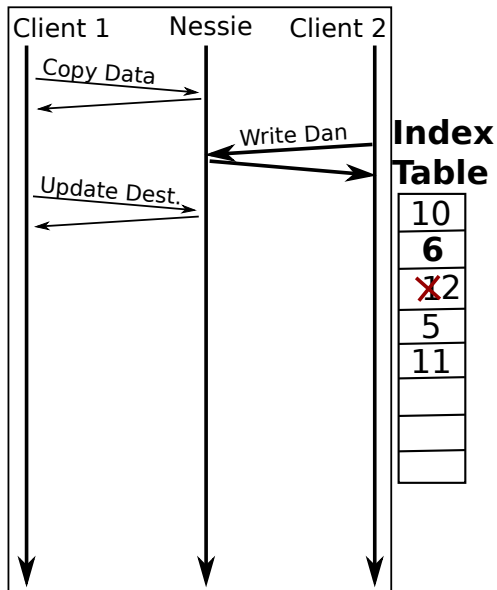
Key	Value	Valid	
Dan	Bacon	T	1
			2
			3
			4
Sam	Pizza	T	5
Dan	Bacon	F	6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



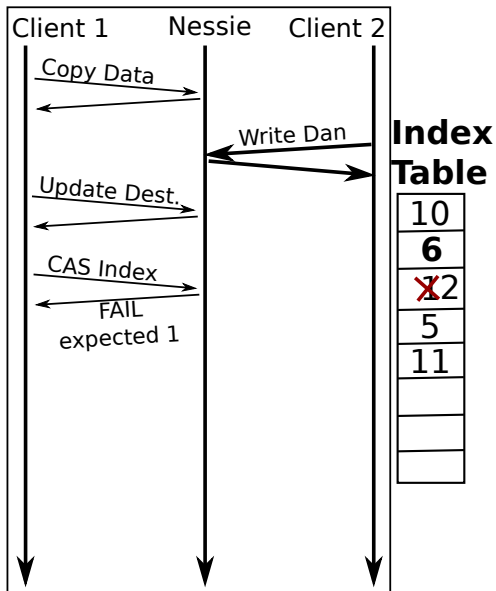
Data Table			
Key	Value	Valid	
Dan	Bacon	T	1
Dan	Tofu	T	2
			3
			4
Sam	Pizza	T	5
Dan	Bacon	F	6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



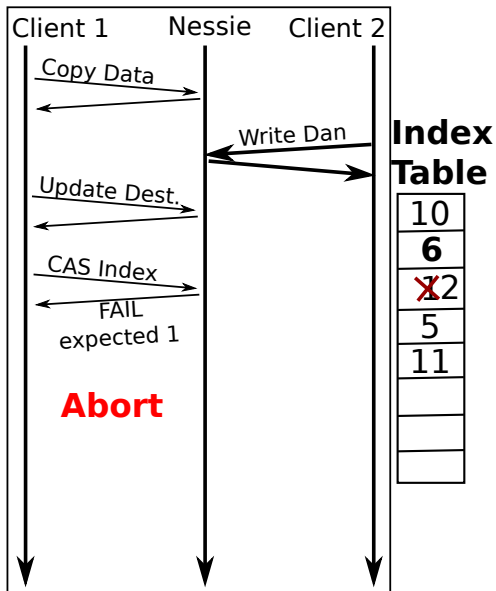
Data Table			
Key	Value	Valid	
Dan	Bacon	T	1
Dan	Tofu	T	2
			3
			4
Sam	Pizza	T	5
Dan	Bacon	F	6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



Data Table			
Key	Value	Valid	
Dan	Bacon	T	1
Dan	Tofu	T	2
			3
			4
Sam	Pizza	T	5
Dan	Bacon	F	6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Nessie: Migration



Data Table			
Key	Value	Valid	
Dan	Bacon	T	1
Dan	Tofu	T	2
			3
			4
Sam	Pizza	T	5
Dan	Bacon	F	6
			7
			8
			9
Ben	Apple	T	10
Ken	Cake	T	11
			12
			13
			14

Performance Tradeoffs

Expected number of roundtrips for a write operation in Nessie:

Load Factor	Roundtrips
0.10	2.5
0.25	3.5
0.50	5.8
0.75	9.9
0.90	13.3

Nessie's Design Complexity

- ▶ Contention over the index table creates complexity
- ▶ Our solution mirrors that of Hardware Transactional Memory
 - ▶ Optimistically attempt operation
 - ▶ Check for conflicting operations
 - ▶ Abort and retry if a conflict is detected
- ▶ NIC support of HTM would greatly improve Nessie
 - ▶ NIC handles detection and reporting of conflicts
 - ▶ Fewer roundtrips

Hardware Transactional Memory

XBEGIN

Read Index Table

Write to Index Table

...

Write to Index Table

XEND

HTM + RDMA

XBEGIN

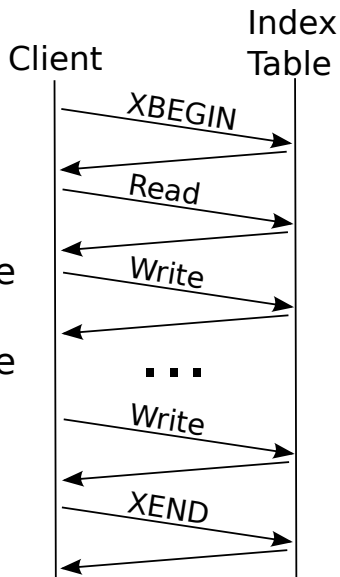
Read Index Table

Write to Index Table

...

Write to Index Table

XEND



Summary

- ▶ Low latency is critical for in-memory distributed systems
- ▶ Nessie:
 - ▶ Exclusively RDMA reads, writes, and CAS operations
 - ▶ Achieves low-latency
 - ▶ Shared lock-free data
- ▶ HTM in conjunction with RDMA can greatly simplify Nessie
- ▶ We are excited to build Nessie and explore HTM