## PYRAMID CODES:

FLEXIBLE SCHEMES TO TRADE SPACE FOR ACCESS EFFICIENCY IN RELIABLE DATA STORAGE SYSTEMS

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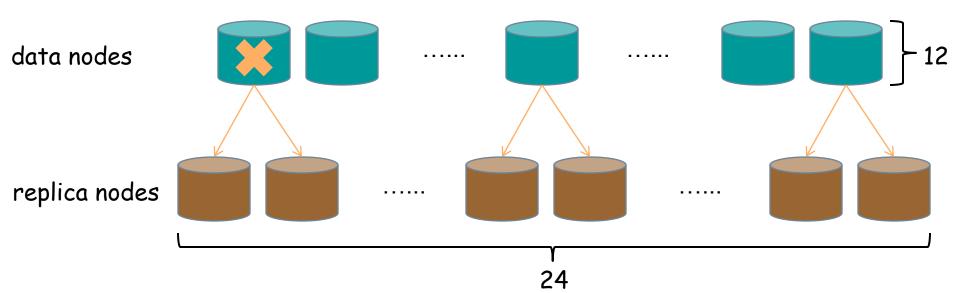
#### networked storage on the rise ...

- rapidly growing demands on storage systems
  - consumers, enterprises ...
  - web services ...
- using commodity components to build large scale
  storage systems is becoming a common practice
  - reliability is a must (five 9's)
  - failure is norm and dealt with by redundancy

#### outline

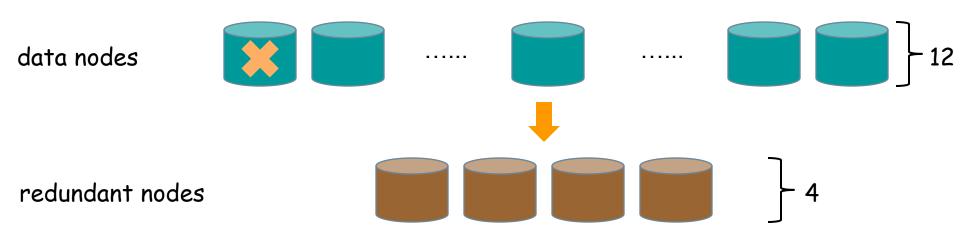
- replication vs. erasure codes
  - the fundamental trade-off
- Pyramid Codes and recoverability theorem
  - not YAC (yet another code)
  - basic Pyramid Codes
  - generalized Pyramid Codes

#### replication vs. erasure codes (1)



- 3-replication
  - storage overhead: 3x
    - 12 data nodes + 24 replica nodes
  - access/recovery cost (one data failure): 1x

#### replication vs. erasure codes (2)



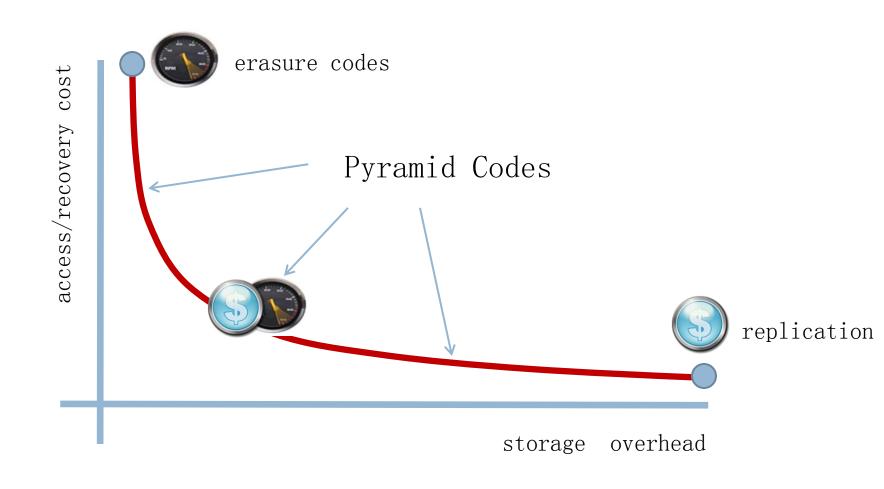
- □ (16, 12) erasure code
  - storage overhead: 1.33x
    - 12 data nodes + 4 redundant nodes
  - access/recovery cost (one data failure): 12x

#### replication vs. erasure codes (3)

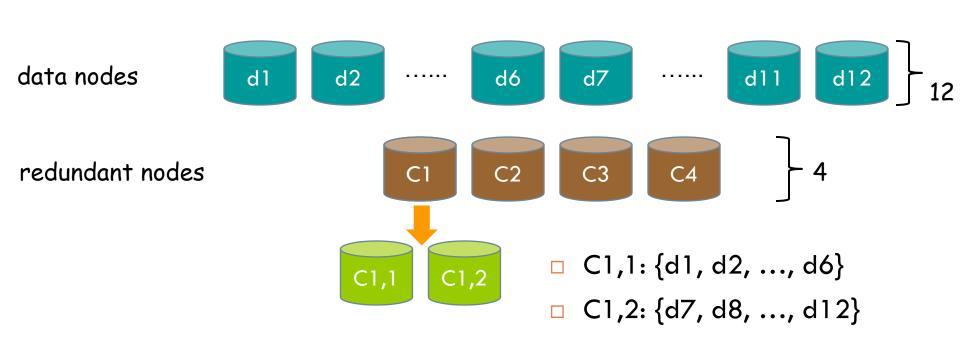
	replication scheme	erasure codes
storage overhead	high (3x)	low (1.33x)
access/recovery cost	low (1x)	high (12x)

- in the end, storage is not that cheap
  - more storage → more machine, more space, more maintenance personal, etc. → 55% of data centers' operating costs (Windows Live service data)
- network traffic is not free either
  - network in data centers can become bottleneck (Lian et al. ICDCS'05)
- □ same concerns for P2P storage ...

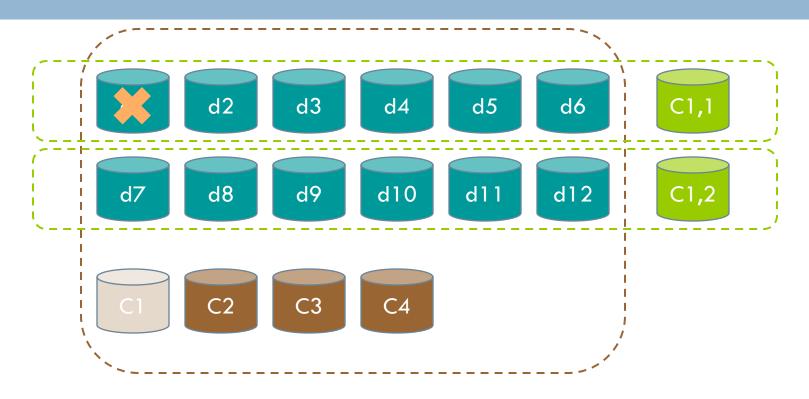
# the fundamental trade-offs in replication vs. erasure codes



#### I. basic Pyramid Codes (1)

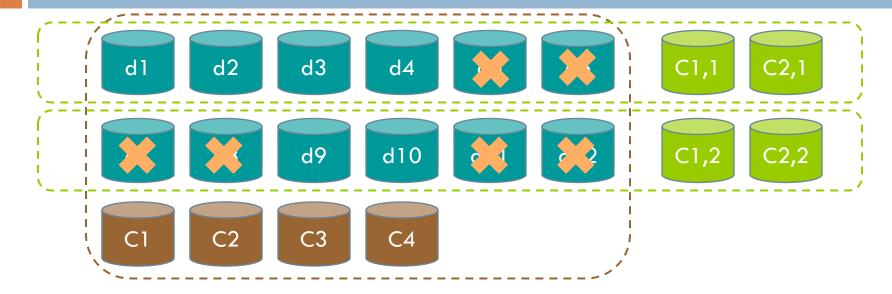


### I. basic Pyramid Codes (2)



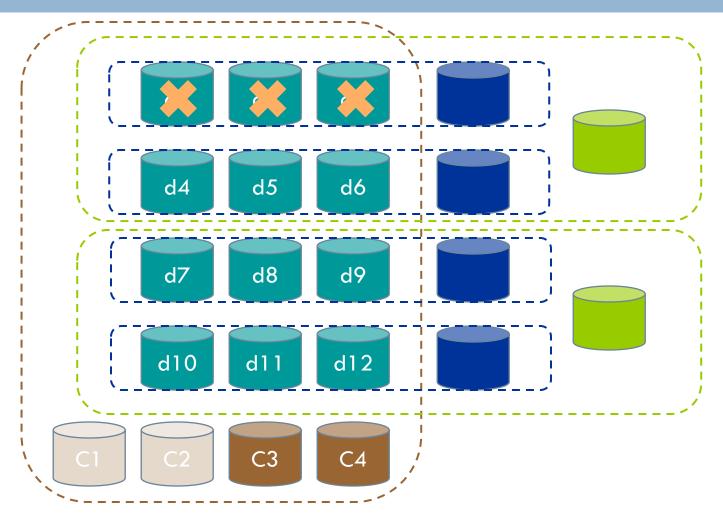
- storage overhead: 1.42x
- □ access/recovery cost (one data failure): 6x
- recovery any 4 failures

#### I. basic Pyramid Codes (3)



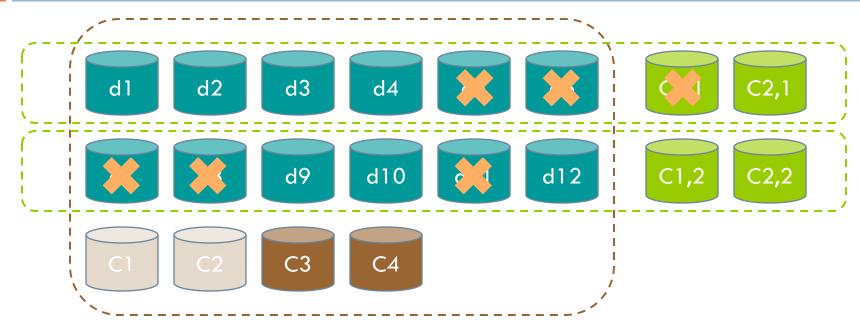
- □ recover d5 and d6
- $\square$  combine C<sub>1,1</sub> and C<sub>1,2</sub>  $\rightarrow$  C<sub>1</sub>; C<sub>2,1</sub> and C<sub>2,2</sub>  $\rightarrow$  C<sub>2</sub>
- recover d7, d8, d11 and d12

#### I. basic Pyramid Codes (4)



decoding is analogous to climbing up a Pyramid!

#### another erasure pattern



- □ is this erasure pattern recoverable at all?
  - no small group recovery

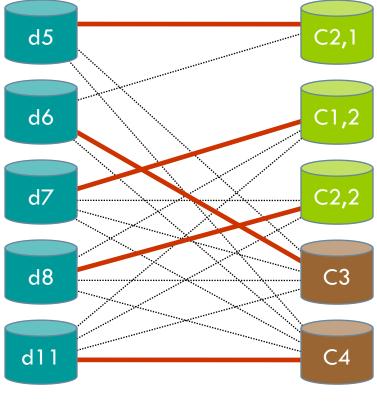
- not recoverable!
- □ C<sub>2,1</sub> and C<sub>2,2</sub>  $\rightarrow$  C<sub>2</sub>, so only 3 redundant nodes at the global level
- counting failures/parities: 5 failures and 5 parities
- □ now what?

#### recoverability theorem (1)

an erasure pattern is recoverable <u>only if</u>
 the corresponding Tanner graph contains

a full-size matching.

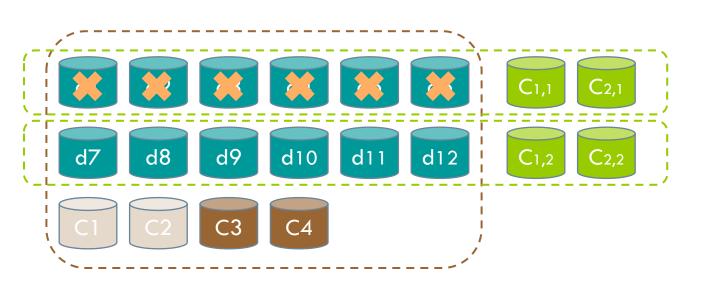
Tanner graph



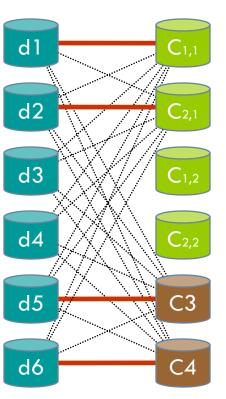
failed data

survival redundancy

### recoverability theorem (2)



an unrecoverable example

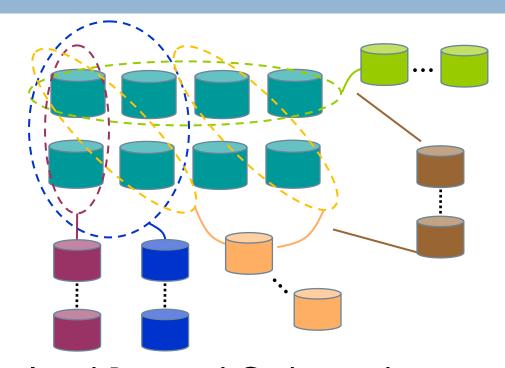


Tanner graph no full-size matching!

### recoverability theorem (3)

- the recoverability theorem is a necessary condition for all erasure codes
- □ it is not sufficient for all known storage codes
  - including basic Pyramid Codes
- generalized Pyramid Codes makes the condition sufficient
  - able to recover any erasure pattern ever possible to recover – optimal recoverably property

### II. generalized Pyramid Codes (1)

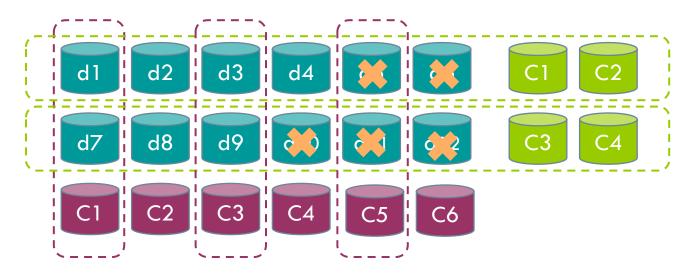


- a generalized Pyramid Code can be constructed given any configuration (data/parity association)
  - details in paper ...
- any generalized Pyramid Code satisfies optimal recoverable property

### II. generalized Pyramid Codes (2)

- why is this a big deal?
  - MDS codes are optimal when redundant nodes and data nodes are fully associated
  - Pyramid Codes are optimal when redundant nodes and data nodes are partially associated
- contributions recap
  - a necessary condition theorem for recoverability
  - a construction algorithm for generalized Pyramid
    Codes, which achieve optimal recoverability

## optimal decoding of generalized Pyramid Codes



- □ how to access/recover with minimum cost?
  - all failed nodes
  - or simply one failed node (say d12)
- optimal decoding path
  - details in paper ...

#### summary

- the fundamental trade-off between
  storage overhead and access/recovery efficiency
- □ two classes of Pyramid Codes
- recoverability theorem
  - generalized Pyramid Codes are optimal