

# Towards Optimal Resource Allocation in Partial-Fault Tolerant Applications

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**Abstract**—We introduce *Zen*, a new resource allocation framework that assigns application components to node clusters to achieve high availability for *partial-fault tolerant* (PFT) applications. These applications have the characteristic that under partial failures, they can still produce useful output though the output quality may be reduced. Thus, the primary goal of resource allocation for PFT applications is to prevent, delay, or minimize the impact of failures on the application output quality. This paper is the first to approach this resource allocation problem from a theoretical perspective, and obtains a series of results regarding component assignments that provide the highest service availability under the constraints imposed by the application data flow graph and the hosting clusters.

We show that (1) even simple versions of this resource allocation problem are NP-Hard, (2) a 2-approximate polynomial-time algorithm works for tree topologies, and (3) a simple greedy component placement performs well in practice for general application topologies. We implement a system prototype to study the application availability achieved by *Zen* compared to failure-oblivious placement, replication, and *Zen*+replication. Our experimental results show that three PFT applications achieve significant data output quality and availability benefits using *Zen*.

## I. INTRODUCTION

With increasing scale and complexity of deployed distributed applications, their utility is increasingly limited by availability rather than performance [10], [29]. Thus, masking failures to ensure application availability has become a key goal in dependable computing. Traditional approaches to fault-tolerance are based on techniques such as replication [4], [12], [15] and checkpointing [12], [13], [18]. Multiple replicas increase the probability that an application will be able to provide service to end-users despite individual node failures. Similarly, checkpointing allows a service to be restarted quickly from a backup site, thereby limiting the work lost when a failure occurs.

However, these approaches introduce well-known tradeoffs between cost and availability. For example, a replicated service may incur significant overheads to provide strict consistency requirements [12], [30]. Further, the monetary cost of implementing highly available services can double for just a fraction of percentage of availability [7], and under correlated failures, even

additional replicas result in a strong diminishing return in availability improvement for many replication schemes [19]. Similarly, the overheads of checkpointing can limit its benefits [31].

In this paper, we investigate an alternative resource allocation framework whose main philosophy is that if we can leverage knowledge of failure characteristics and resource capacity constraints (e.g., CPU, MEM, I/O, etc.) of node clusters in *component placement* (i.e., to which cluster should each component be assigned), we can achieve significant gains in availability for an important class of applications, which we term *partial-fault tolerant* (PFT) applications. In contrast to applications that require the availability of all components to operate correctly, PFT applications provide a “graceful degradation” in performance as the number of failures increases. For example, aggregation systems such as Sawzall [21], SDIMS [28], and PIER [11] are likely to be able to tolerate some missing objects while processing a query (e.g., JOIN, MEDIAN, AVG, etc.) on a distributed database. Similarly, data mining applications such as WTTW [26] and FAB [25] (described in Section II-A) can still classify data objects under failures, though with less confidence. Further, for many data stream applications with stringent temporal requirements [2], [3], it is more important to produce partial results within a given time bound than full results produced late. Finally, mission-critical applications (e.g., power control, telecom, medical systems) deploy multiple sensors [8] such that at least some of them should be able to trigger an alert during failures or when operating conditions are violated (e.g., when a sensor senses a fire.)

Given this system model, we can precisely state the problem as follows: Given a distributed computing system comprising  $n$  clusters ( $T_1, T_2, \dots, T_n$ ) each with a resource capacity  $c_i$  and a failure probability  $p_i$  ( $i \in [1, n]$ ), and a PFT application made up of  $m$  components ( $C_1, C_2, \dots, C_m$ ) each of which may execute on any cluster, allocate each of the  $m$  modules to one of the  $n$  clusters such that the *loss in expected application*

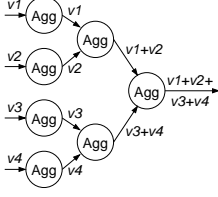


Fig. 1: Data aggregation application.

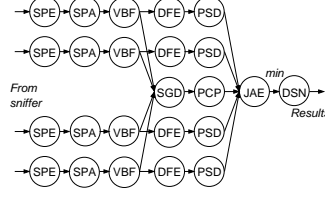


Fig. 2: WTTW data flow graph.

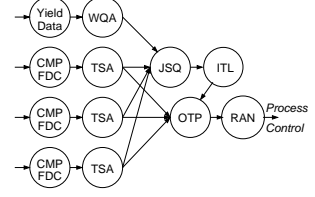


Fig. 3: FAB data flow graph.

output value is minimized under failures subject to the constraints imposed by the application data flow graph, the resource capacities, and the failure probabilities.

To solve this problem, we introduce the Zen framework that performs *failure-aware component assignment* to achieve high availability for PFT applications. Zen approaches the problem from a theoretical perspective by reducing the above optimization problem to a variant of the graph coloring problem. We show that even simple versions of this graph coloring problem are NP-Hard. Therefore, Zen provides (1) a 2-approximate (wrt the optimal algorithm's output loss) polynomial-time algorithm for tree topologies and (2) a greedy algorithm using the underlying principles of the optimal algorithm for general graph topologies, and shows that it works well in practice.

We have implemented a prototype of Zen and evaluated its effectiveness by deploying across a cluster testbed three real-world PFT applications: a data aggregation application, WTTW, and FAB. Relative to a failure-oblivious placement, Zen improves availability of these applications by 35%-60%. Compared to replication, (1) Zen provides nearly equal availability as replication but with less cost and (2) Zen+replication achieves significantly higher availability from 33%-90% for the same cost as replication.

By using Zen, we hope that application managers can significantly improve service availability by making an *application-aware* and *cluster-aware* automated resource allocation. Our long-term goal is to enable self-adaptive component placement for achieving high availability in large-scale, dynamic environments. Towards this goal, this paper makes three technical contributions:

- We introduce the problem of failure-aware component assignment for PFT applications and show that even simple versions of this problem are NP-Hard.
- We develop the Zen resource allocation framework that provides a 2-approximate algorithm for tree topologies and a greedy algorithm for general application topologies that works well in practice.
- We implement a prototype of Zen in a large-scale stream processing system and evaluate its effect on availability compared to failure-oblivious placement,

replication, and Zen+replication. Our experimental results show that three PFT applications achieve significant availability benefits using Zen.

The rest of this paper is structured as follows. In Section II, we discuss motivating applications and examples. Section III defines our system model in terms of the class of failures and applications we target. Section IV gives a proof of the NP-hardness and describes our approximation algorithms. Section V describes evaluation of Zen using our system prototype. Finally, Section VI discusses related work and Section VII concludes.

## II. MOTIVATION

### A. Motivating Applications

Our work is motivated by PFT applications. In this section, we describe three PFT applications that span a broad range of size, topology, functionality, and fault-tolerance. For each application, we focus on quantifying its availability in terms of the application output quality.

**Data Aggregation.** Our first application is a data aggregation service that is designed to be representative of in-network query processing systems such as Sawzall [21] and SDIMS [28]. In our implementation, an aggregation tree is constructed (Figure 1) that assigns aggregator components at internal nodes which receive children inputs and send the aggregate output to their parent. An aggregator's data output value is well-expressed as a summary of its incoming data values. Several standard aggregation functions such as SUM, AVG, MAX, MIN, etc. satisfy this application model. Availability in this application is measured by the percentage of paths from the leaves to the root that contribute to the global aggregate. E.g., a COUNT query can tolerate some missing data inputs to count the number of nodes in a network.

**WTTW.** The Who's Talking to Whom (WTTW) application is a VoIP-based stream-processing application [26]. Its aim is to identify and track conversing parties in real-time from distributed and noisy/compressed speech signals. The data flow graph of WTTW (Figure 2) shows the inputs from various remote sniffers, each carrying several compressed speech signals, that feed data to different operators for processing and classifying conversation pairs. The details of these operators are beyond the scope of this paper and are described in [26].

In WTTW, all components' output values are a SUM of their inputs except JAE that performs a JOIN operation. WTTW's quality metric (described in Section V) is defined in terms of the accuracy in correctly identifying conversation pairs [26].

**FAB.** The third PFT application, FAB, monitors processing of silicon wafers in a semiconductor manufacturing environment [25]. Its goal is to predict wafer yield from summary vectors input from the tool sensors and detect any tool anomalies in real-time. FAB does self-learning as it continuously evaluates its prediction performance against the ground truth i.e., actual wafer yield. Figure 3 shows data flow graph of FAB; details of FAB operators are described in [25]. FAB's quality metric (described in Section V) is defined as its prediction accuracy that depends on the path availability from the sensors to the OTP classifier and the incremental learner ITL.

### B. Motivating example

**A simple yet subtle example.** Consider the PFT application in Figure 4: component 1 computes a SUM (say) over outputs of components 2 and 3. Figure 4 shows three possible component assignments:  $\alpha$ ,  $\beta$ , and  $\gamma$ . In ( $\alpha$ ), we assign root ( $C_1$ ) to one cluster (black) and  $C_2$  and  $C_3$  to another cluster (gray). In ( $\beta$ ), we assign  $C_1$  and  $C_3$  to the gray cluster and  $C_2$  to the black cluster, and in ( $\gamma$ ), we assign  $C_1$ ,  $C_2$ , and  $C_3$  to the gray cluster.

Note that allocation  $\beta$  is better<sup>1</sup> than  $\alpha$  because if the black cluster fails, then the application output for  $\alpha$  goes to 0. However, under allocation  $\beta$ , the system could still process data flowing from  $C_3$  to  $C_1$ . Indeed, if the gray cluster fails, both allocations give no output. A careful calculation<sup>2</sup> shows that the best allocation, however, is  $\gamma$  that keeps all components on the same cluster. The main intuition behind this is that only one cluster failure affects allocation  $\gamma$  while two cluster failures can hinder the other allocations.

There are several important observations from this example. First, allocate as many components as possible to the same cluster (subject to cluster resource constraints) to maximize expected output value under failures. Second, assign components on independent paths to different clusters to avoid dependent failures. Finally, for heterogeneous clusters with different failure probabilities, assign "highly important" components to clusters with the lowest failure probabilities. We use these observations in designing our component placement algorithms in Section IV.

<sup>1</sup> Assuming a uniform cluster failure probability  $p$ , the failure probabilities (i.e., no output from the root) are  $FP(\alpha) = p + (1-p)p$ ,  $FP(\beta) = p$ , and  $FP(\gamma) = p$ ;  $FP(\beta) = FP(\gamma) \leq FP(\alpha)$ .

<sup>2</sup> The expected output values are  $E[V]_\gamma = (V_a + V_b)(1-p)$  and  $E[V]_\beta = V_b(1-p)p + (V_a + V_b)(1-p)^2$ ;  $E[V]_\gamma \geq E[V]_\beta$ .

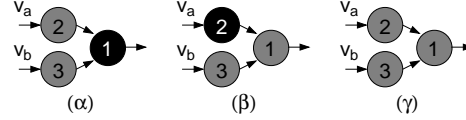


Fig. 4: Three possible assignments of components to 2 clusters.

## III. SYSTEM MODEL

In this section, we present the failure and application models based on which we formulate the optimization problem of failure-aware component allocation.

### A. Failure model

While many previous studies assume failures to be independent, in reality, the assumption of failure independence is rarely true. Node failures are often correlated, with multiple nodes in the system failing nearly simultaneously [9], [19], [22], [23]. E.g., nodes on a rack connected by a switch or a power supply, network-attached disks, nodes running the same version of a vulnerable software [15], etc. Further, the size of correlated failures can be quite large and have a dominant affect on system availability [19]. We therefore use the model proposed by Junqueira et al. [16] which groups all resources that can fail together into clusters where each cluster can fail-stop in its entirety, causing all its resources and hosted application components to be unavailable. Note that a single cluster failure also affects other components that depend on the output from its failed components. Individual cluster failures are considered independent and identically distributed. Each cluster<sup>3</sup>  $T_i$  has a failure probability  $p_i$ .

### B. Application model

A PFT application topology is modeled as a directed graph of inter-connected components where vertices are either data sources, sinks, or processing operators, and edges represent the data flow between them. We assume each input source  $s$  (e.g., a leaf in a tree) having a certain *importance value*  $v_s$  (scalar value) defined as  $s$ 's contribution to the application output. Alternatively,  $v_s$  is the "loss" incurred in the application's total output value if  $s$  fails. The importance value of a component  $C$  (output) is expressed as a linear combination<sup>4</sup> or a MIN/MAX of its input importance values. Intuitively, each  $C$  performs "value addition" to its data inputs that can be modeled as a linear function in many applications. E.g., for a tree topology,  $v_i$  of an internal node  $i$  is conceptually a SUM of  $v_l$  for all leaves  $l$  lying in the subtree rooted at  $i$ . Further, this simple model fits well the real applications described in Section II-A.

<sup>3</sup> Nodes failing independently can be mapped to clusters of size one.

<sup>4</sup>  $\alpha_{e_1} v_1 + \alpha_{e_2} v_2 + \dots + \alpha_{e_k} v_k$  where  $\alpha_{e_i}$  is the weight associated with edge  $e_i$  receiving input  $v_i$  from the  $i$ th child of  $C$ .

For more complex applications, we require application designers to specify the relative importance values of the components. The importance metric allows us to rank application components so that highly important components get assigned to highly reliable clusters.

#### IV. PROBLEM HARDNESS AND ALGORITHMS

First, we show that even simple cases of the component assignment problem (Section I) are NP-Hard. Second, for applications with tree topologies, we provide a polynomial-time optimal algorithm assuming an unbounded number of homogeneous clusters ( $c_i = c, p_i = p \forall i$ ), and a 2-approximate (wrt optimal's output loss) algorithm under a bounded number of clusters. Third, for general graph topologies, we present a greedy algorithm using the underlying principles of the optimal algorithm and show that it works well in practice in Section V.

##### A. Optimal Component Allocation is NP-Hard

To prove optimal component allocation is NP-Hard, consider a simple tree-based PFT application (e.g., Figure 1) computing an aggregation function. For this application, we first characterize the effect of cluster failures on the output value using the following lemma:

*Lemma 4.1:* Let  $X : P \rightarrow N$  be some fixed placement of components to clusters. For a leaf  $l$ , let  $P(l)$  be the set of components on the path from  $l$  to the root, and let  $S(l)$  denote the set of clusters that have at least one component in  $P(l)$  assigned to them. Let  $v_l$  denote the importance value of the input entering  $l$ . If  $p_j$  denotes the probability of failure of cluster  $j$ , then the expected output value of the application is

$$\sum_{l \in \text{leaves}} (\prod_{j \in S(l)} (1 - p_j)) v_l$$

If all the failure probabilities are equal i.e.,  $p_j = p \forall j$ , then the expected output value becomes

$$\sum_{l \in \text{leaves}} (1 - p)^{d(l)} v_l$$

where  $d(l) = |S(l)|$  is the number of distinct clusters on the path  $P(l)$ .

*Proof:* Let us consider input data that enters a leaf  $l$  and follows the unique path  $P(l)$  from  $l$  to the root. If no cluster (and hence component) on this path fails, then this stream contributes a value  $v_l$  to the root aggregate value. Otherwise, if any component fails on  $P(l)$ , this input contributes 0. The probability that no component on the path  $P(l)$  fails is exactly the probability that no cluster in  $S(l)$  fails, equal to  $\text{pr}(l) = \prod_{j \in S(l)} (1 - p_j)$  since cluster failures are independent. Thus the expected contribution of leaf  $l$  to the root output is  $\text{pr}(l) * v_l$ . By

linearity of expectation, the expected root output value is  $\sum_{l \in \text{leaves}} \text{pr}(l) v_l$ . If  $p_j = p \forall j$ , then  $\prod_{j \in S(l)} (1 - p_j) = (1 - p)^{|S(l)|} = (1 - p)^{d(l)}$  by definition of  $d(l)$ . ■

Given Lemma 4.1, we can view the component assignment problem as the following tree coloring problem: Given a rooted tree  $T = (V, E)$  with vertex set  $V$  and edge set  $E$ , assigning components optimally to a set of  $k$  clusters of capacity  $c_j (j \in [1, k])$  is equivalent to coloring  $V$  with colors  $1, 2, \dots, k$  such that color  $j$  is used no more than  $c_j$  times and the term  $\sum_{l \in \text{leaves}} (1 - p)^{d(l)} v_l$  is maximized, where  $d(l)$  is the number of distinct colors on the path from leaf  $l$  to the root.

Intuitively, the goal is to color the tree vertices such that each path has as few colors as possible. Translating back, we want to minimize the total number of distinct clusters used on every path from a leaf to the root. Observe that if no cluster fails, the maximum achievable output value is  $\sum_{l \in \text{leaves}} v_l$ . Thus, we can define the *loss* of a placement to be the difference between the maximum achievable value and the value achieved by the placement i.e.,  $\text{loss} = \sum_{l \in \text{leaves}} v_l (1 - (1 - p)^{d(l)})$ . Hence, maximizing the expected output value is equivalent to minimizing the loss<sup>5</sup>. Taking the first order approximation, our objective function of minimizing the loss in the output value becomes:

$$\text{Min} \sum_{l \in \text{leaves}} (1 - (1 - p)^{d(l)}) v_l \approx \text{Min} \sum_{l \in \text{leaves}} p d(l) v_l$$

Now, we can define the loss minimization problem as:

*Definition 4.2 (Loss minimization problem):* Given a tree  $T(V, E)$  with input values  $v_l$  at leaf  $l$ , color  $V$  from colors  $1, 2, \dots, k$  such that color  $i$  is used at most  $c_i$  times and the term  $\sum_l v_l d(l)$  is minimized where  $d(l)$  is the number of distinct colors on the path from  $l$  to root.

Notice that  $p$  has been dropped from the loss definition since it is a fixed constant.

*Lemma 4.3:* The loss minimization problem is NP-Hard even for tree instances with maximum degree 3 and all colors (clusters) have equal capacities.

*Proof:* We sketch the proof by showing a reduction from the 3-partition problem: Given  $3m$  non-negative numbers  $a_1, \dots, a_{3m}$ , is there a partition of these numbers into  $m$  sets such that the sum of each set is equal to  $S/3$  where  $S = \sum_{i=1}^{3m} a_i$ ? This problem is NP-Hard even when  $S$  is polynomially large in  $m$ . The main idea of the reduction is as follows: Consider a tree with  $3m$  disjoint paths of length  $a_1, a_2, \dots, a_{3m}$ , all connected to a common root, and  $m$  clusters each with capacity  $S/3$ . Since our goal is to minimize the number of distinct

<sup>5</sup>However in terms of approximation guarantees, the two quantities are different. E.g., if an optimal solution achieves 99% of the maximum achievable value, then a 2-approximate algorithm wrt value will only guarantee at least 49.5% value, while a 2-approximation wrt loss will guarantee at least 98% value. Zen provides 2-approximation wrt loss.



clusters on each path, it can be shown that there exists a placement in which each path lies in a single cluster iff the 3-Partition problem has a solution [5]. ■

It is important to note that the reason for NP-hardness does *not* directly result from coloring paths with fewest colors. Rather, it arises due to the limited number of colors available as we next show a polynomial-time optimal algorithm assuming an unbounded number of colors<sup>6</sup> and adapt this solution to obtain a 2-approximate algorithm when the number of colors is bounded.

### B. Optimal Component Allocation Algorithm

First, we present a simple lemma that provides a useful structural property for designing the optimal tree coloring algorithm. Consider an optimum solution OPT. We can assume that OPT satisfies the following property without loss of generality.

*Lemma 4.4 (Consecutive Vertex Color Property):*

Consider a vertex  $v$  and another vertex  $v'$  such that  $v'$  lies in the subtree rooted at  $v$ . If  $v$  and  $v'$  are both assigned color  $Q$ , then every component on the path between  $v$  and  $v'$  must also be assigned the color  $Q$ .

*Proof:* We will show that any solution can be modified in such a way that this property is satisfied without worsening the value of the solution. Consider a solution that does not satisfy this property. Then, there must exist a vertex  $v$  and another vertex  $v'$  in the subtree rooted at  $v$  such that both have color  $Q$ , but the child of  $v$  does not have color  $Q$ . Call this vertex  $v''$  and let its color be  $Q''$ . Now, consider the solution obtained by swapping the colors of  $v'$  and  $v''$ . Note that this swap can only possibly affect the value  $d(l)$  (and hence the contribution to the overall solution) of the leaves that lie in the subtree rooted at  $v''$ . Let  $L'$  be the set of leaves that lie in subtree rooted at  $v'$  (and hence also  $v''$ ) and let  $L''$  denote the leaves that lie in the subtree rooted at  $v''$  but do not lie in the subtree rooted at  $v'$ . Note that for leaves in  $L'$ , the colors along the path to root do not change, hence  $d(l)$  remains unchanged for these leaves. For the leaves in  $L''$ , the color  $Q''$  is replaced by  $Q$ , however since  $v$  was already colored  $Q$ , and as each leaf to root path for leaves in  $L''$  also passes through  $v$ , the quantity  $d(l)$  can only decrease for each such leaf. Thus the result follows. ■

**Algorithm sketch.** Consider a tree  $T$  under OPT color assignment. By Lemma 4.4, all nodes colored  $i$  form a connected subgraph. Let  $r'$  be a node colored  $i$  and let  $T'$  denote the subtree rooted at  $r'$ . Consider the trees  $T'$  and  $T'' = T \setminus T'$ . Since we want to color paths with the least

number of colors, only color  $i$  can be shared<sup>7</sup> between  $T'$  and  $T''$ , and hence the only connection between the solutions for  $T'$  and  $T''$ . We use this observation to compute  $V(T', q)$  starting bottom-up from the leaves where  $V(T', q)$  denotes the optimum solution for  $T'$  under the constraint that the color of  $T'$ 's root is used at most  $q$  times. Clearly, computing OPT corresponds to computing  $V(T, c)$  where any color can be used at most  $c$  times.

Next, we show how to compute  $V(T', q)$  bottom-up: Suppose we color the root  $r$  by color 0. Let  $T_1, \dots, T_k$  be the subtrees rooted at children of  $r$ . Using Lemma 4.4, we know that for a subtree  $T_i$ , either its root is also colored 0, or the color 0 is never used in the subtree  $T_i$ .

**Case 1:** If a subtree  $T_i$  does not use color 0, we claim that the loss that  $T_i$  contributes to  $T$  is  $\text{loss}(T_i) + \sum_{l \in L_i} v_l$ . Here  $L_i$  is the set of leaves in  $T_i$  and  $\text{loss}(T_i)$  is the minimum achievable loss for  $T_i$ . This holds because the subtree  $T_i$  does not contain the color 0, but the parent of the subtree,  $r$ , is colored 0. Hence for every path from a leaf in  $T_i$  to the root, the value of  $d(l)$  (i.e., the number of colors from a leaf to the root) increases by exactly 1. Since the second term ( $\sum_{l \in L_i} v_l$ ) is a constant, it suffices to find the best solution to  $T_i$  subject to the constraint that it does not contain color 0.

**Case 2:** If the subtree  $T_i$  uses color 0, the root of  $T_i$  must be colored 0. For every leaf in  $L_i$ , since each path to the root of  $T_i$  already has color 0, the number of distinct leaf to root colors  $d(l)$  in  $T$  is equal to the number of distinct colors from the leaf to the root of  $T_i$ . Thus, the loss contribution of  $T_i$  to  $T$  is  $\text{loss}(T_i)$ . Hence it suffices to compute the best solution to  $T_i$  subject to the constraint that the root of  $T_i$  is colored 0.

We use these observations to design a polynomial-time dynamic programming based optimal algorithm [5] for the loss minimization problem. The optimal solution runs in time  $O(|V|^2 c^2)$  and space  $O(|V|^2 c)$ .

### C. 2-Approximate Algorithm

We next show a 2-approximate algorithm wrt loss to OPT under a bounded number of resource clusters.

*Lemma 4.5:* There is a 2-approximation with a bounded number of colors (clusters.)

*Proof:* Compute the optimal solution as described above using unbounded resources. Note that no color is used more than  $c$  times in this solution. Moreover this solution is clearly a lower bound on the optimum achievable loss assuming bounded resources. Let  $m_i$  denote the number of times color  $i$  is used in the unbounded

<sup>6</sup>Note that unbounded number of colors doesn't make the problem trivial since we still need to color the paths with least number of colors.

<sup>7</sup>Having an unbounded number of colors enables this sharing since we are not constrained by the *total* capacity.

solution. We have  $\sum_i m_i = m$ , where  $m = |V|$  and each  $m_i \leq c$ ; note that  $m \leq c * n$  must hold for a feasible solution. We line the  $m_i$ -colored nodes up together one after another, and form blocks of size  $c$ . Note that since each  $m_i \leq c$ , any set  $m_i$ -colored nodes can lie in at most two blocks. Each block corresponds to a color in the new solution. In any leaf to root path in the new solution, the number of distinct colors used is at most twice the number of colors used in the original unbounded solution. Thus the result follows. ■

#### D. Greedy Component Allocation Algorithm

Finally, we present a greedy component placement for general graph topologies under heterogeneous clusters and components.

Using the importance metric (Section III-B), the loss minimization term  $\sum_{l \in \text{leaves}} v(l)d_l$  (Section IV-A) can be equivalently expressed as  $\sum_C I(C)Z(C)$  where  $I(C)$  is the importance of component  $C$ , and  $Z(C)$  is 0 if  $C$  and its parent are placed on the same cluster and 1 otherwise [5]. Thus, to minimize loss, set  $Z(C) = 0$  for as many adjacent components as possible, especially the ones with the higher values of  $I(C)$ .

This formulation suggests two guiding principles consistent with our observations in Section II-B: (1) components of higher importance should be placed on clusters with higher capacities (and low failure probabilities) i.e., choose cluster  $j$  with highest  $\frac{c_j}{p_j}$  and (2) all components lying on a path from a source to the sink should be co-located on the same cluster (if possible) i.e., minimize the total number of colors (clusters) on all paths.

Given an application data flow graph  $G(V, E)$ , Algorithm 1 allocates components in decreasing importance to clusters ranked by  $\frac{c_j}{p_j}$  ( $j \in [1, n]$ ). It defines a connected subgraph  $SG$  of components that are co-located on the same cluster (say  $T$ ) as follows: at each step, assign the highest importance  $C_k$  to  $T$  (if spare capacity) if  $(C_k, SG) \in E$  i.e.,  $(C_k, C_p) \in E$  for some  $C_p \in SG$ .

#### V. EXPERIMENTAL EVALUATION

In this section, we evaluate Zen's availability benefits for three real-world PFT applications (Section II-A) compared to failure-oblivious placement (FOP), replication, and replication+Zen. For replication, we create one replica for the highest importance components of each application and assign components to clusters randomly. FOP uses a simplified version of our topology-aware approach by co-locating components in a path on the same cluster but without considering failure probabilities.

##### A. Methodology

The Zen prototype and the hosted PFT applications run across a cluster testbed on top of the Stream Pro-

#### Algorithm 1 Greedy Component Placement Algorithm

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1: Calculate  $I(C)$  for components  $C = \{C_1, C_2, \dots, C_m\}$ .
2: Rank the clusters  $T_1, T_2, \dots, T_n$  sorted (decreasing) on  $\frac{c_j}{p_j}$ 
   ( $j \in [1, n]$ )
3:  $j \leftarrow 1$ 
4: while  $C \neq \emptyset$  do
5:   Select the highest importance component  $C_i \in C$ 
6:   while  $T_j$  has spare capacity do
7:     Assign  $T_j \leftarrow C_i$ ;  $C \leftarrow C \setminus \{C_i\}$ ;  $SG \leftarrow \{C_i\}$ 
8:     Select highest importance  $C_k \in C$  s.t.  $(C_k, SG) \in E$ 
9:     if  $\exists C_k$  satisfying (8 :) AND  $T_j$  has capacity then
10:       $T_j \leftarrow C_k$ ;  $C \leftarrow C \setminus \{C_k\}$ ;  $SG \leftarrow SG \cup \{C_k\}$ 
11:     else {no such  $C_k$  exists OR  $T_j$  has no capacity left}
12:       break;
13:     end if
14:   end while
15:   if  $T_j$  has no spare capacity then
16:      $j \leftarrow j + 1$ 
17:   end if
18: end while

```

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cessing Core [14]. The cluster testbed is an IBM BladeCenter comprising 100 nodes with Intel Xeon and AMD Opteron processors with 2 GB to 4 GB RAM running Linux kernel 2.6.21 connected by 1 Gbit Ethernet. The BladeCenter is arranged into networked chassis where each chassis has up to 13 blades (nodes) and denotes a dependent failure unit (Section III-A) that can fail due to switch misconfigurations, power failures, etc.

**Failure Model.** Based on prior studies, we use two failure models: (1) FM1 representing a controlled stable enterprise system and (2) FM2 a dynamic wide-area network e.g., PlanetLab [1]. Both models classify nodes into three classes with availabilities of 0.3, 0.8, and 0.99 but differ in the class size i.e., number of nodes per class.

- 1) FM1 is based on a one year failure study of a 400-node cluster that observed about 4% nodes accounted for 70% failures, 25% nodes accounted for 20% failures, and 70% nodes accounted for 1% failures [22]. Correspondingly, we approximate class sizes of 5%, 25% and 70%, respectively.
  - 2) FM2 is based on a three month failure study of 240 PlanetLab nodes that showed about 6% nodes had less than 30% availability, 50% nodes had about 80% availability, and 37.7% nodes had more than 99% availability [6]. Correspondingly, we approximate class sizes of 10%, 50%, and 40%, respectively.
- We expect a typical computing environment to exhibit failure properties between FM1 and FM2 and thus, our results will quantify best and worst-case performance.

##### B. Data Aggregation Application Results

We implemented an aggregation tree comprising 63 components or processing elements (PEs) [14] on 6

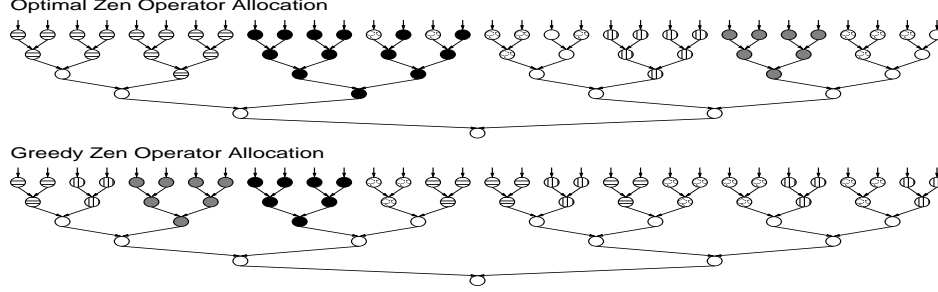


Fig. 5: Optimal Zen and Greedy Zen placement for the data aggregation tree. Each pattern represents one chassis.

chassis with each PE running on a single host. At each leaf, a source continuously generates data packets and each internal node computes the SUM aggregate of the number of updates in its subtree. We quantify this application’s quality metric as the ratio of root’s output value to the total data input through all the leaves.

Figure 5 shows the placement for optimal Zen and greedy Zen for this application. Note that both optimal and greedy Zen aim to satisfy the consecutive vertex color property—the average number of colors per path is 2 in both cases. The difference between their allocations is because of the most important PEs in the tree (the ones closest to the root.) Greedy Zen places PEs close to the root on one failure unit (chassis); all nodes with depth 0, 1, 2 share the same chassis. In contrast, the optimal algorithm includes paths from two leaves to the root so that the application can produce output even if all the clusters but the one marked white fail.

Figures 6(a) and 6(b) show the effect of optimal Zen, greedy Zen, replication, and optimal Zen+replication on the quality metric by varying the number of failed clusters for FM1 and FM2, respectively. For replication, nodes at levels 0 (root), 1, and 2 are replicated on different machines. Each bar value represents the expected quality, and the following results are computed relative to optimal Zen. For FM1, optimal Zen increases quality by 10%-17% compared to greedy Zen, and 37%-72% compared to replication. For FM2, the corresponding numbers are 7%-11% and 33%-73%, respectively. Since replication only focuses on important components but does not optimize global allocation, it does not yield high benefits under failures that might be a limiting factor for resource-constrained applications. As the number of failed clusters increase, both optimal and greedy Zen show a graceful degradation in the output quality. Finally, Zen+replication outperforms all allocations.

### C. WTTW Results

Figure 7 shows the assignments for 39 WTTW PEs using greedy Zen and FOP on 8 chassis. Note that Zen assigns the highest importance PEs, DSN, JAE, PCP, and

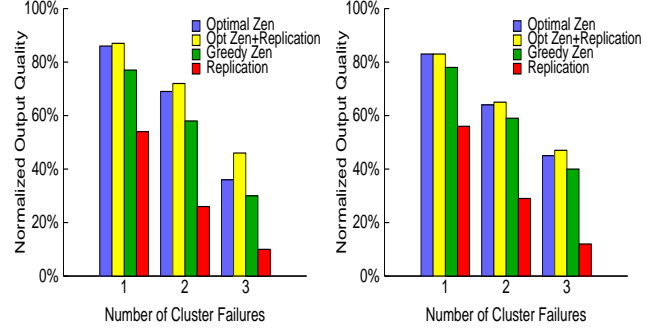


Fig. 6: Comparing optimal Zen, optimal Zen+replication, greedy Zen, and replication for the data aggregation application under failure models FM1(a) and FM2(b).

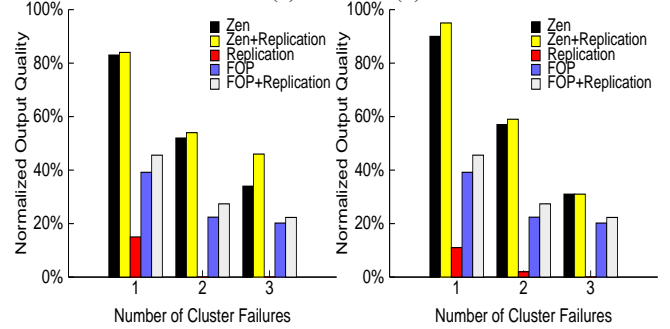


Fig. 8: Normalized  $Q_{WTTW}$  for Zen, Zen+replication, replication, FOP, and FOP+replication under FM1(a) and FM2(b).

SGD, on the highest reliability chassis A. Since FOP is failure-oblivious, it assigns them to a less reliable chassis F. The WTTW’s quality metric  $Q_{WTTW}$  is defined in terms of the accuracy in correctly identifying conversation pairs [26] as:

$$Q_{WTTW} = \frac{(1 - \text{DSN error rate})(\text{No. of pairs reported})}{(\text{No. of conversation pairs})}$$

where DSN is the output PE in Figure 2. We collected results from 600 input conversation streams.

Figures 8(a) and 8(b) show the effect of Zen, Zen+replication, replication, FOP, and FOP+replication on  $Q_{WTTW}$  (normalized wrt no-failure quality) for different number of cluster failures under FM1 and FM2, respectively. Each bar value is computed as an average of 8 independent runs, and all numbers below are repre-

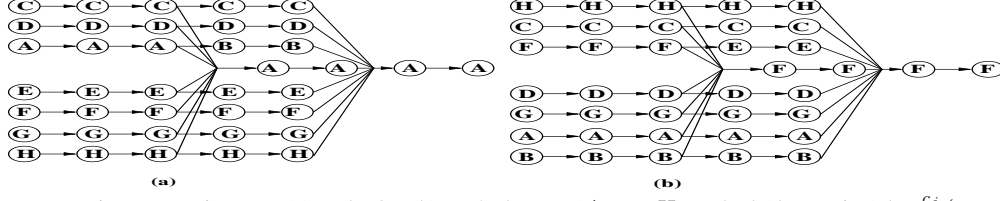


Fig. 7: WTTW PE assignment using Zen (a) and FOP (b) on 8 chasses ( $A, \dots, H$ ) ranked (decreasing) by  $\frac{c_j}{p_j}$  ( $j \in \{A, \dots, H\}$ ).

sented relative to Zen. For FM1, Zen increases quality by 80% compared to replication, 40%-50% compared to FOP, and 35%-45% compared to FOP+replication. For FM2, the corresponding numbers are 90%, 35%-60%, and 30%-50% respectively. The Zen+replication strategy outperforms all other strategies, achieving up to 30% quality increase over Zen.

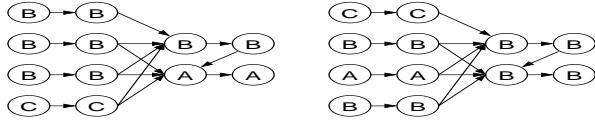


Fig. 9: FAB PE Assignment using Zen (a) and FOP (b) on 3 chasses ( $A, B, C$ ) ranked (decreasing) by  $\frac{c_j}{p_j}$  ( $j \in \{A, B, C\}$ ).

#### D. FAB results

The third application, FAB, comprises 12 PEs. Compared to data aggregation and WTTW, FAB is (1) small-scale and (2) expected to operate primarily in local area networks that exhibit failure properties closer to the FM1 model. Therefore, we use two FM1-based failure models for FAB: FM1a comprising three failure classes with availabilities 75%, 90%, and 99%, and FM1b with 90%, 95%, and 99%. Our findings, unavoidably, depend on the failure models we used. However, we believe these models are representative for actual deployments of the application.

Figure 9 shows FAB's PE assignments using greedy Zen and FOP on 3 chasses  $A, B, C$  ranked (decreasing) by  $\frac{c_j}{p_j}$  ( $j \in \{A, B, C\}$ ), and having 2, 8, 4 nodes, respectively. Note that Zen places the two most important PEs, OTP and RAN, on the most reliable chassis  $A$ . FOP does topology-aware but failure-oblivious assignment assigning these PEs to chassis  $B$ . For replication, OTP and RAN replicas are placed together on one cluster. The FAB quality metric  $Q_{FAB}$  at time  $t$  is defined in terms of the accuracy in yield output [25]:

$$Q_{FAB}(t) = \frac{\hat{N}(t)}{N(t)} [P_D(t) - 0.2P_F(t)]$$

where  $N(t)$  is the total number of wafers,  $\hat{N}(t)$  the number of wafers processed by OTP,  $P_D$  and  $P_F$  are the probabilities of fault-detection and false positives, respectively. We ran FAB on real data from 9000 wafers, and present results averaged across time over 1000 runs.

Figures 10(a) and 10(b) show the effect of all the five component allocations on  $Q_{FAB}$  (normalized wrt quality under zero failures) for different number of cluster failures under FM1a and FM2b, respectively. The following results are computed relative to Zen. For FM1a, Zen increases quality by about 43% compared to replication, 17%-20% compared to FOP, and 3%-13% compared to FOP+replication. For FM1b, the corresponding numbers are 37%-45%, 4%-18%, and nearly the same quality as FOP+replication, respectively. As in previous cases, Zen+replication outperforms all strategies, and increases quality by up to 12% over Zen.

In summary, our evaluation shows that for the three real-world PFT applications we considered, Zen's component allocation can significantly reduce the loss in an application output quality under failures.

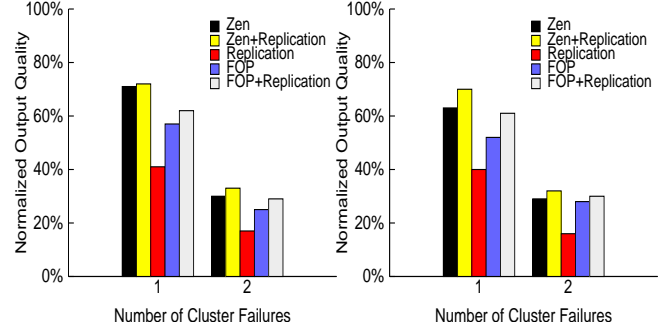


Fig. 10: Normalized  $Q_{FAB}$  for Zen, FOP, replication, and Zen+replication under failure models FM1a (a) and FM1b (b).

#### VI. RELATED WORK

Several studies [6], [9], [22], [23], [31] (and the references therein) have aimed at characterizing the failure properties of real systems. We leverage two important observations common among these studies in our failure model: (1) node failures are not uniformly distributed, and a small fraction of nodes incur most of the failures [6], [22], [31] and (2) node failures are often correlated, with multiple nodes in the system failing nearly simultaneously [9], [19], [22], [23]. Some of these studies have also used knowledge of failure characteristics in resource allocation to improve cluster availability [9], [31].

Similarly, fault-tolerance techniques have been extensively studied for improving availability. E.g., replica-



tion, erasure coding, placing replicas on nodes running heterogeneous software versions [15], etc. in distributed storage systems. In task allocation systems, previous solutions [17] have aimed at maximizing the probability of running the *entire* task successfully but not for the PFT model. To our knowledge, there has not been prior work on task allocation for PFT applications.

In stream processing systems, component (operator) placement techniques have generated recent interest primarily for improving performance in resource-limited sensor networks [24] and wide-area stream systems [20], [27]. For achieving high availability, research in this area has focused on replication [4], and storing data and checkpointing [12] but not on failure-aware component allocation. These techniques are complementary to Zen in that it can leverage them to further enhance availability as shown in Section V.

Finally, two of the design principles in Zen are closely related to the observations by Yu et al. in analyzing availability of multi-object operations [29]: (1) concentrate objects on fewer machines for “strict” operations, and (2) spread objects across machines for “more tolerant” operations. A strict operation can be viewed as a path from a source to a sink requiring availability of *all* in-between nodes, and thus we co-locate them on as few clusters as possible. A more tolerant operation corresponds to multiple parallel paths that should be assigned to different clusters for fault-tolerance.

## VII. CONCLUSIONS AND FUTURE WORK

We introduced Zen, a new failure-aware resource allocation framework to achieve high availability for PFT applications, and show that even simple versions of the component placement problem are NP-Hard. Therefore, Zen provides a 2-approximate polynomial-time algorithm for tree topologies and a greedy algorithm for general graph topologies. Our evaluation shows that for three real-world PFT applications, Zen significantly reduces loss in the application’s output quality under failures.

Our future work is to enable self-adaptive component placement to improve availability in large-scale, dynamic environments. Further, we plan to integrate both failure-aware and network-aware (proximity) properties in Zen to achieve optimal component assignment that simultaneously provides high availability and high performance.

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