

Abstraction-Driven Network Verification and Design (a personal odyssey)

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It started in 2004

- A sabbatical at CMU
 - Joined a collaborative project with AT&T Labs
 - Goal: To reverse engineer the routing designs of 100s of production networks and find ways to detect errors early and minimize outages due to routing loops and blackholes
 - We were given only router configuration files

Excerpts of a Router Config

```
interface Ethernet0
ip address 6.2.5.14 255.255.255.128
interface Serial1/0.5 point-to-point
ip address 6.2.2.85 255.255.255.252
ip access-group 143 in
frame-relay interface-dlci 28
```

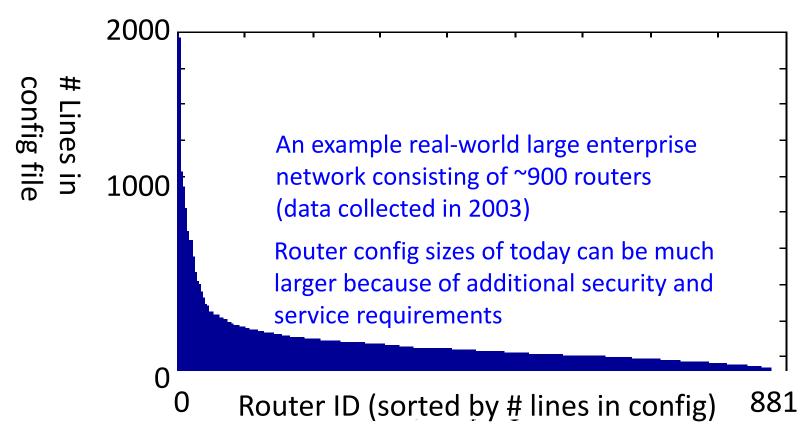
router ospf 64
redistribute connected subnets
redistribute bgp 64780 metric 1 subnets
network 66.251.75.128 0.0.0.127 area 0
router bgp 64780

access-list 143 deny 1.1.0.0/16 access-list 143 permit any route-map 8aTzlvBrbaW deny 10 match ip address 4 route-map 8aTzlvBrbaW permit 20 match ip address 7 ip route 10.2.2.1/16 10.2.1.7

redistribute ospf 64 match route-map 8aTzlvBrbaW neighbor 66.253.160.68 remote-as 12762 neighbor 66.253.160.68 distribute-list 4 in

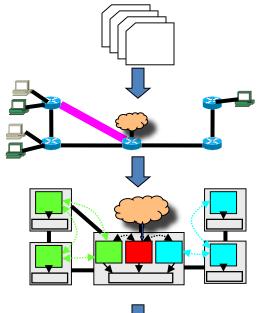
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Lots of Configuration Files



A Reverse-Engineering Methodology

[Maltz et al, Sigcomm'04]



BGP AS1

Configuration files
Find links

Construct logical IP Topology

Find adjacent routing processes

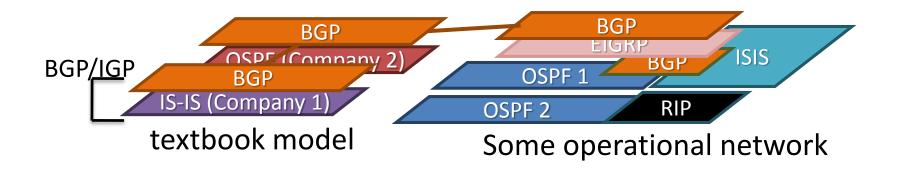
Construct Routing Process Graph

Condense adjacent routing processes into

Routing Instances
This abstraction is key to a scalable solution

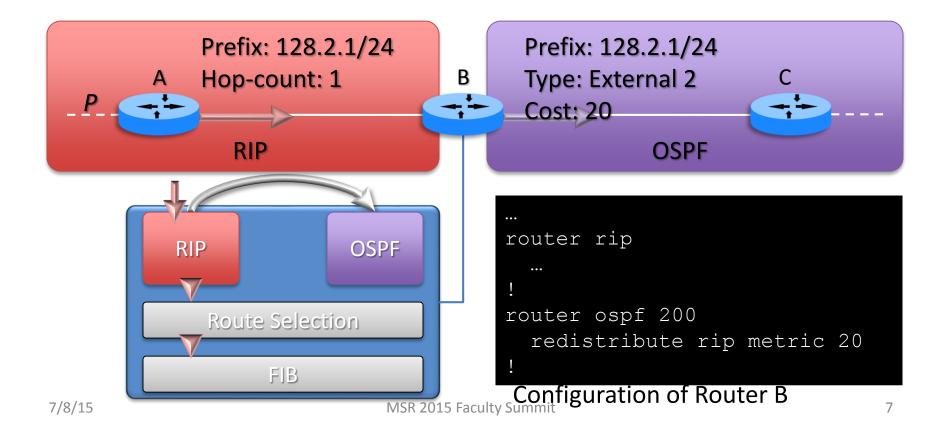
OSPF #1

Diverse Deigns beyond Textbook Model

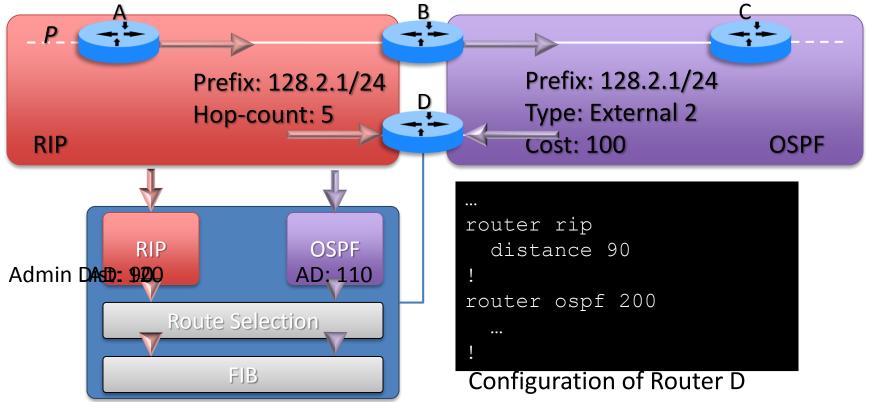


A network can have many routing instances and their interaction will impact routing safety

Route Redistribution



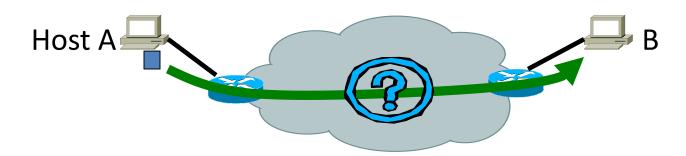
Route Selection



Impact on Routing Theory & Practice

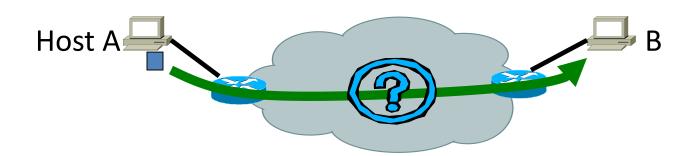
- Understanding current design and ensuring safety
 [Le et al, ICNP'07, Sigcomm'08, CoNext'10] [Benson et al,, IMC'09, NSDI'09] [Alim and Griffin, CoNext'11] [Sun et al, CoNext'12], etc.
- Clean slate design of route redistribution [Le et al, Sigcomm'10]
- Routing Reconfiguration
 [Vanbever et al, Sigcomm'12] [Vissicchio et al, Infocom'14] etc.
- Co-existence of multiple control planes (including SDN) [Volpano et al, HotSDN'14] [Vissicchio et al, Infocom'15] etc.

Next goal: Predicting Reachability

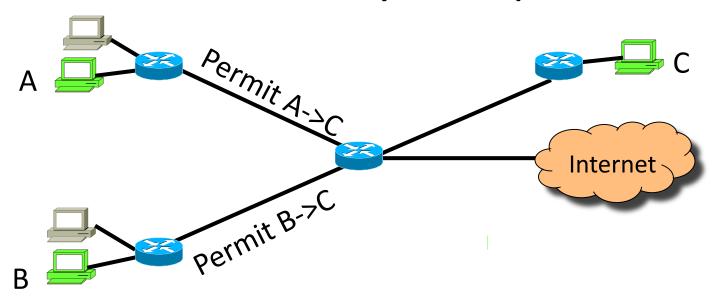


- Reachability depends not just on topology
 - Routing protocols, packet filters, and middleboxes
- Predicting reachability is key to network security and resilience

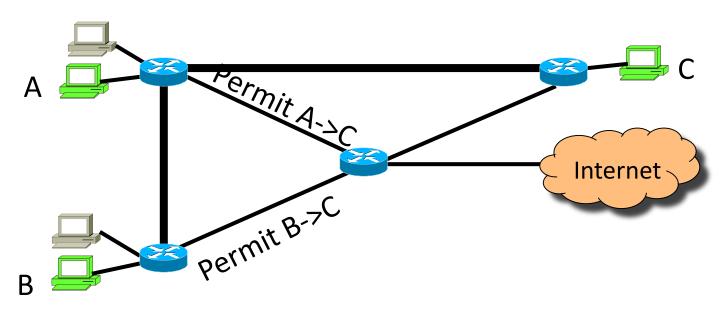
State of the Art at the time



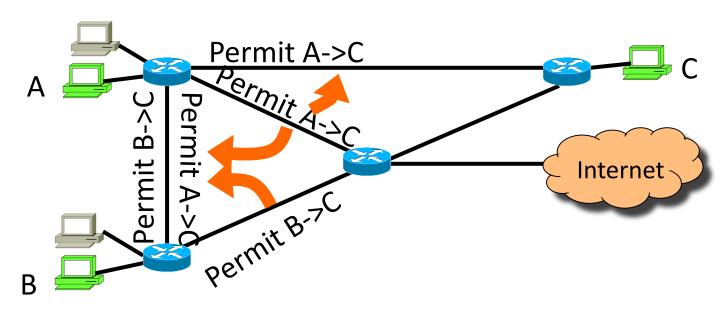
- Build the network and try it
- Dynamic probing (ping and traceroute) used to troubleshoot reachability problems



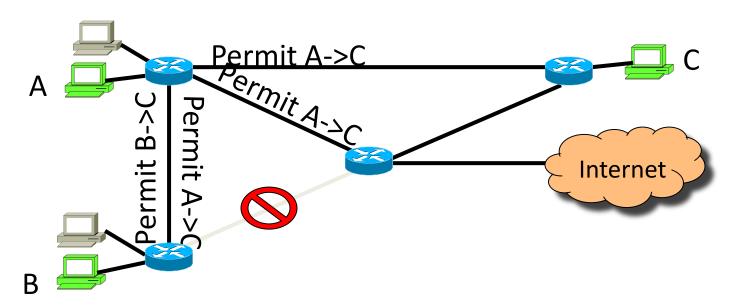
- Enterprise with two remote offices
- Only A&B should be able to talk to server C

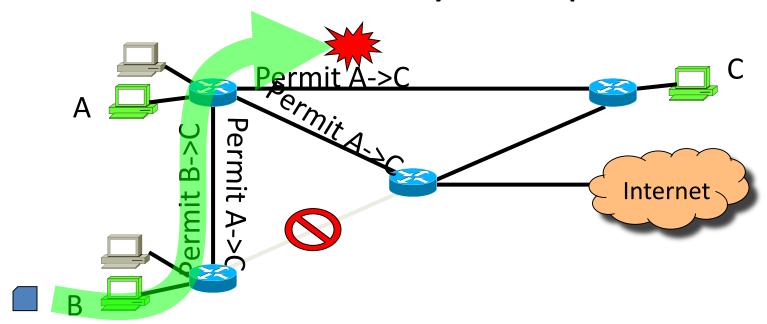


- Network designers add two links for robustness
- Configure routing protocols to use new links in failure



Designers apply packet filters to new links





- Packets from B->C dropped!
- Testing under normal conditions won't find this error!

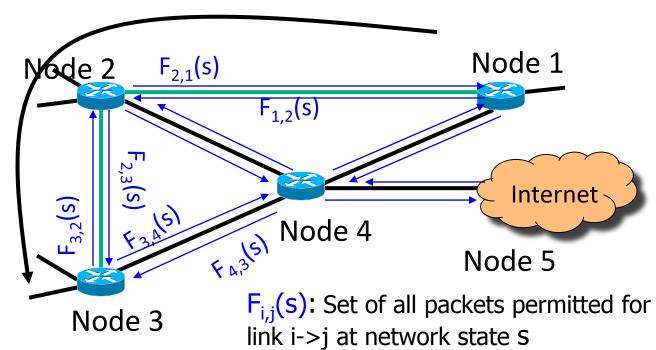
The Reachability Set abstraction

[Xie et al, Infocom'05]

Set of all packets permitted from one node to another

- Model packet filters naturally
 e.g., "Permit A->C" rule defined on link from node u to v:
 F_{u,v} = {packet p | p.src_addr = A, p.dst_addr = C}
- Effect of routing protocols added as <u>dynamic</u> destination address based packet filters
 - when network is in forwarding state s, $F_{u,v}(s) = F_{u,v} \cap \{\text{all packets } u \text{ would forward to } v \text{ at } s\}$
- Packet transformation as generalized inverse function that maps a set of packets to another set of packets

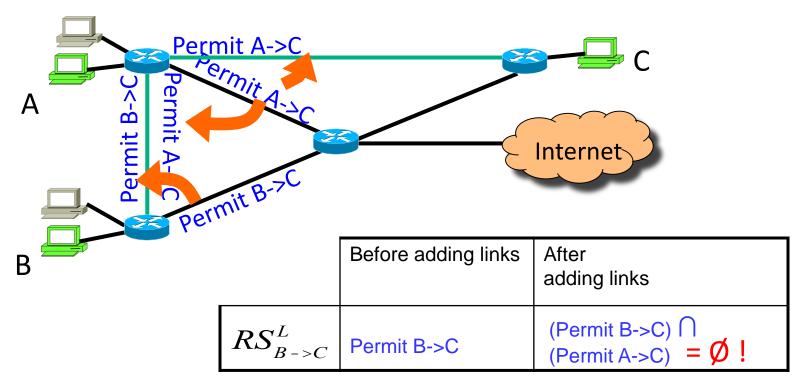
Reachability Analysis Graph



Reachable Set over directed path $1->2->3=F_{1,2}(s)\cap F_{2,3}(s)$

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Let's revisit that reachability example



Recent Advances in Static Network Analysis

- Boolean satisfiability formulation
 [Mai et al, Sigcomm'11]
- Header space analysis [Kazemian et al, Sigcomm'12]
- Fast algorithms [Yang and Lam, ICNP'13]

Next goal: Design automation

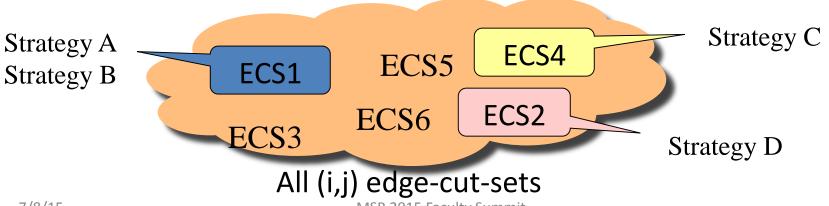
[Sung et al CoNext'10]

- 1. Abstract network-wide requirements of a design task
 - Correctness criteria for reachability control modeled as a Reachability Matrix: each cell $\mathbf{RS}(i,j)$ defines precisely the required reachability set from (virtual) subnet i to j
- 2. Formulate optimization problems
 - Incorporate resource feasibility constraints
 (e.g., router capacity for processing packet filter rules)
 - Model explicitly operator strategies (e.g., to deploy a minimum number of filter rules)
- 3. Solve formulated problems
 - Obtain new packet filter placement algorithm

Automated Packet Filter Placement

Intuition:

- To achieve RS(i, j), same filters must be replicated in an edge-cut-set (ECS) between gateways of the subnets
 - Correctness guaranteed
- Variety of heuristics possible based on design strategy which chooses particular ECS (minimizing total # of filters, balancing processing load, etc.)



Some Related Design Efforts

- Integrated design methodology [Sun and Xie, CoNext'13]
- Optimizing the "one big switch" abstraction [Kang et al, CoNext'13]
- Placement of middleboxes and NFVs
 [Anwer et al, SOSR'15]

Conclusion

- A huge <u>semantic gap</u> exists between network service objectives and actions of individual protocols and nodes.
 - Software defined networking (SDN) doesn't reduce service objectives, while introducing a new type of nodes and diverse control apps
- Developing <u>higher level</u> abstractions may be key to containing this "curse of many knobs".
 - E.g., Separation of correctness and performance concerns as in traditional computer programming?